



Parallel Computation Of 2D Morse-Smale Complexes

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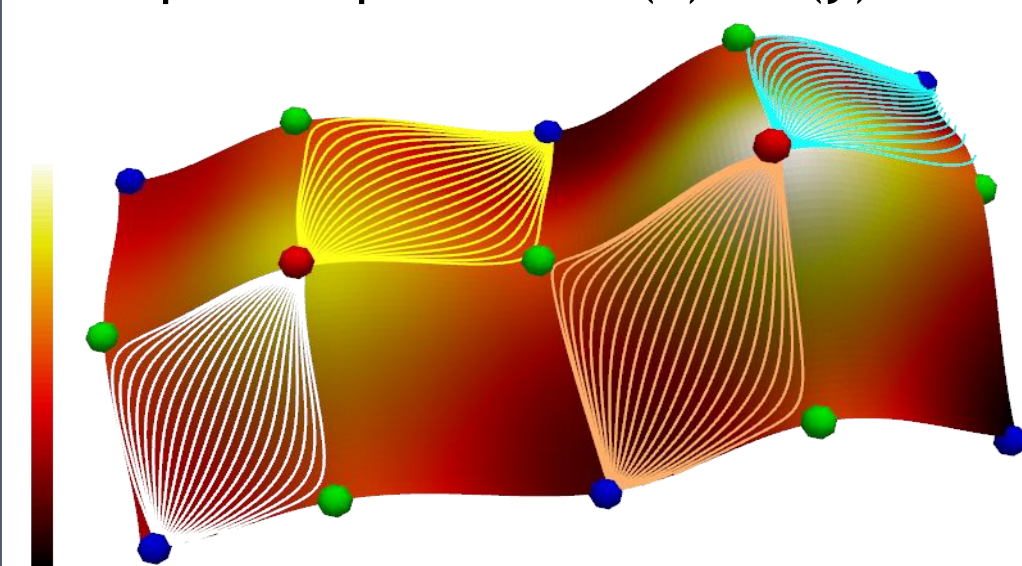
Objective

Parallel computation of the Morse-Smale complex designed for multi-core and GPU environments.

Computation for datasets that don't fit in Memory.

The Morse-Smale Complex

Input: Scalar (real valued) function
Example: 2D plane $\rightarrow \sin(x) + \sin(y)$.



Gradient curves: Curves that trace the direction of steepest descent

Critical Points: Points of origin/destination of gradient curves

Morse-Smale complex: Partition based on origin/destination of gradient curves.

Reference

Shivashankar, N. Maadaswamy, S. and Natarajan, V., Parallel Computation of 2D Morse-Smale Complexes, *IEEE Transactions on Visualization and Computer Graphics* (To Appear).

Contact and Acknowledgements

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Work supported by Intel and DST.

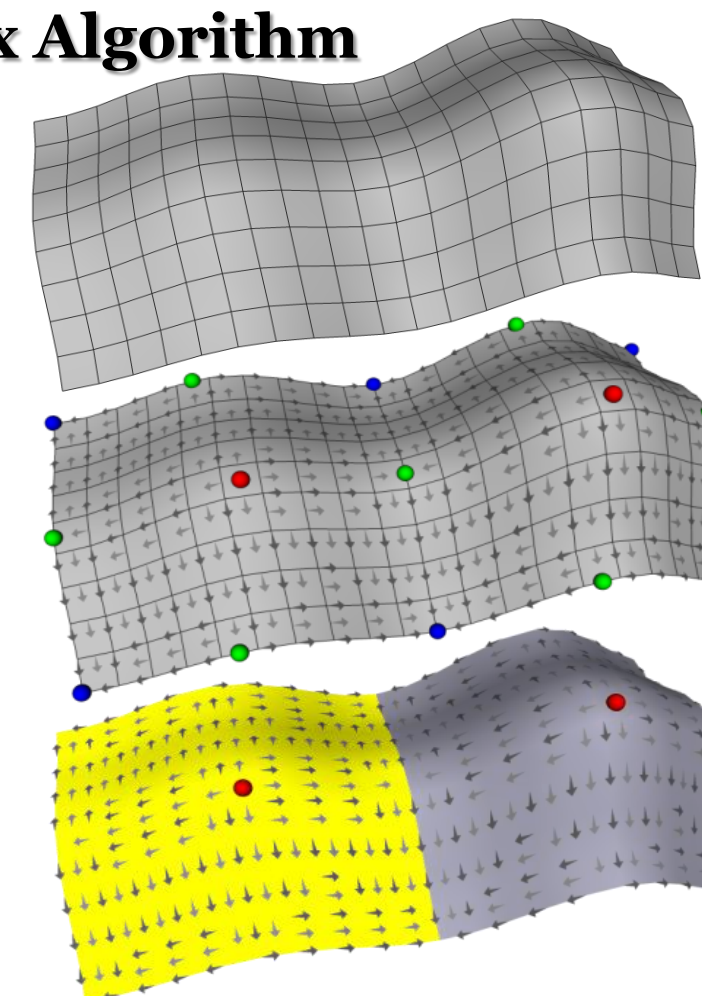
Parallel Morse-Smale(MS) Complex Algorithm

INPUT: Domain as a Cell complex and a scalar function sampled at vertices.

1. Compute discrete gradient

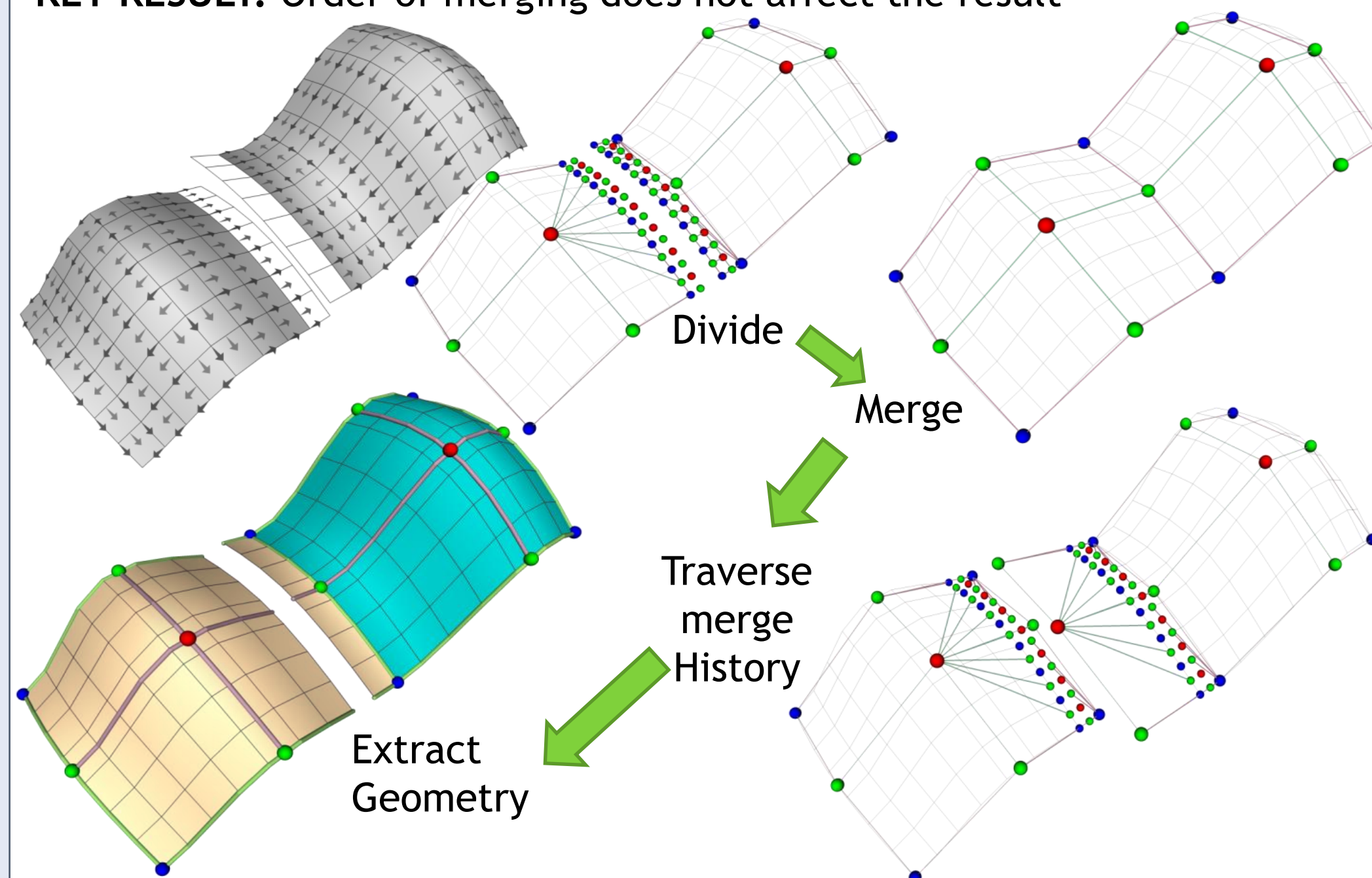
KEY RESULT: Computing discrete gradient is independent of order. Can be done in parallel.

2. Traverse the gradient field to compute the MS complex.



Large Data: Divide and conquer.

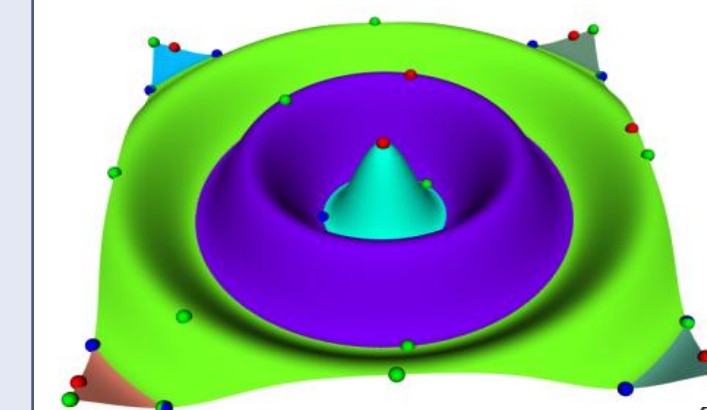
KEY RESULT: Order of merging does not affect the result



Evaluation

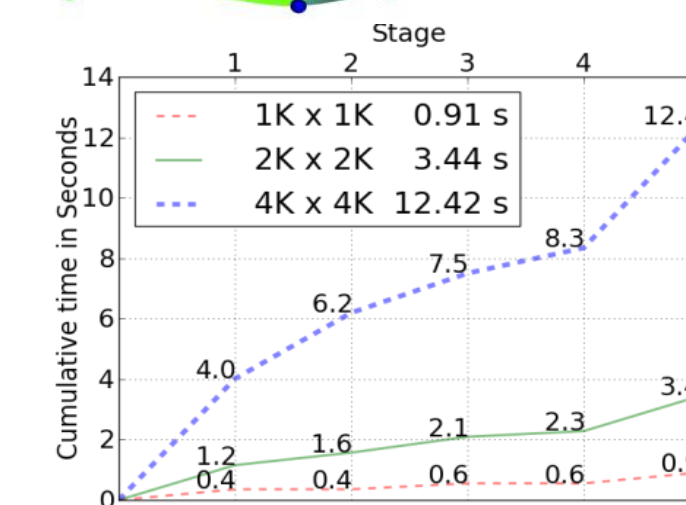
Experiments with synthetic and real world datasets on 2D grids.

Observed near linear scaling with data size and number of cores.

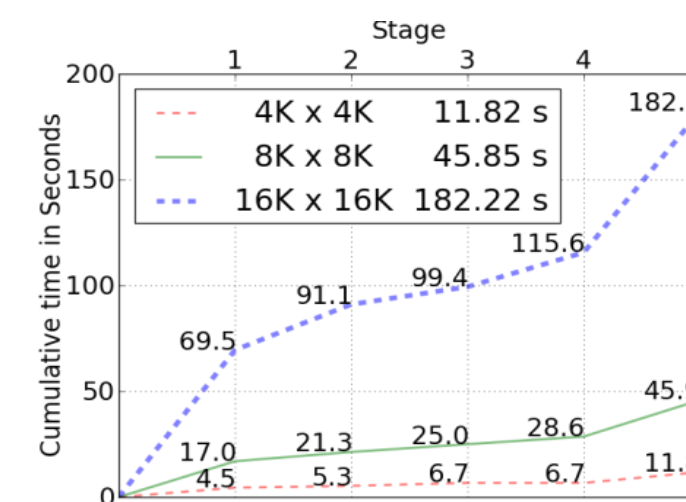


WGAUSS dataset

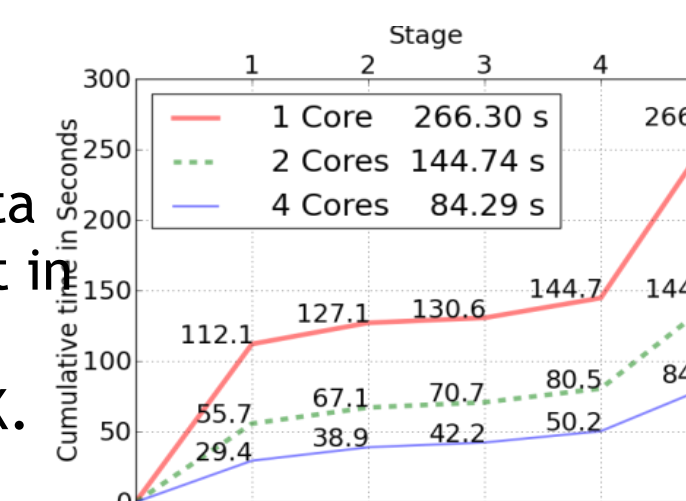
Data fits in CPU but not GPU



Data fits neither in CPU nor GPU.



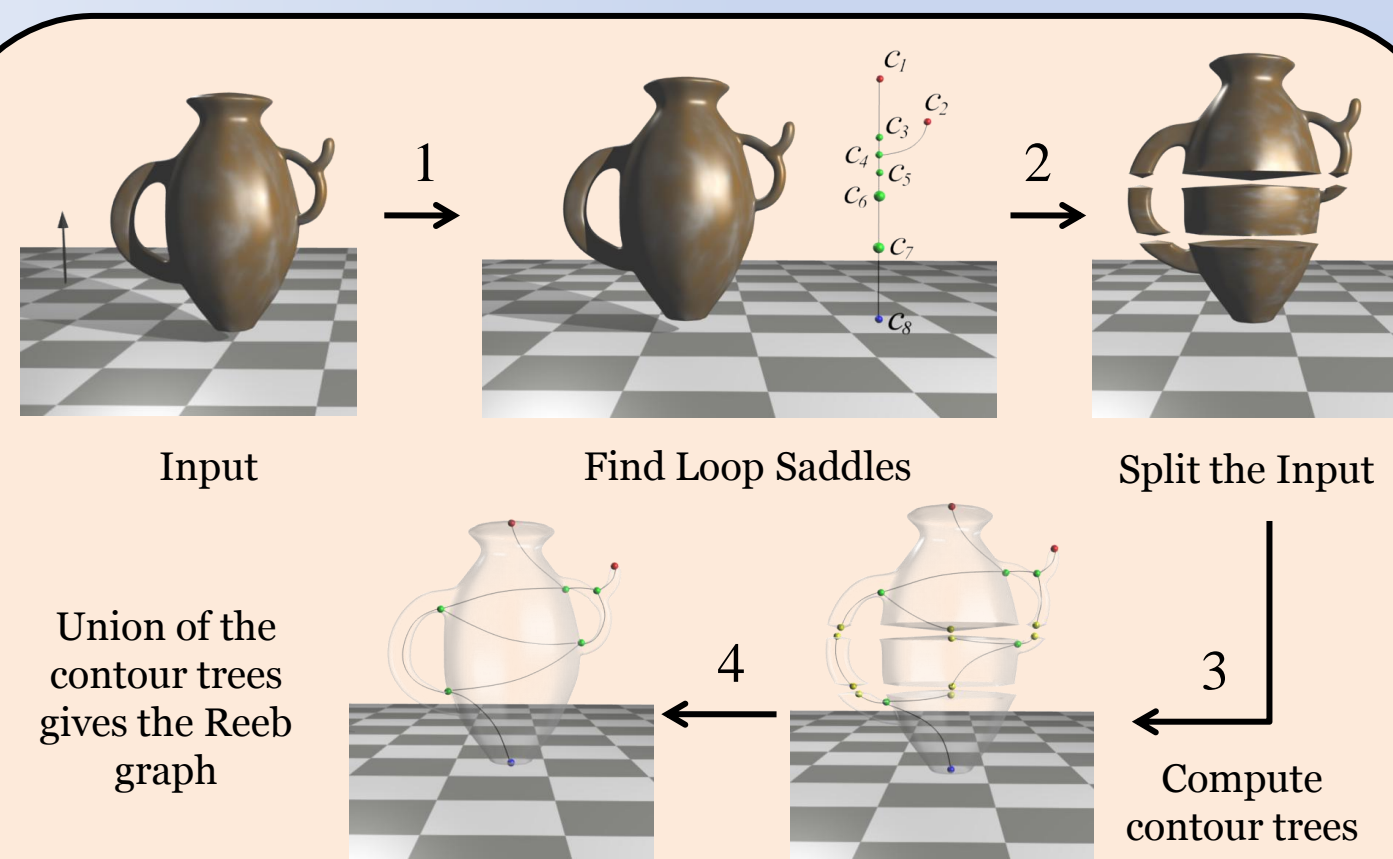
Computed in CPU. Data does not fit in CPU.
Size: 8Kx8K.



The Reeb graph of a scalar function is obtained by mapping each connected component of its level sets to a point

Applications

- Topology based shape matching. *Hilaga et al. SIGGRAPH 2001*
- Transfer function design. *Weber et al. TVCG 2007*
- User Interface. *Bajaj et al. Vis 1997*



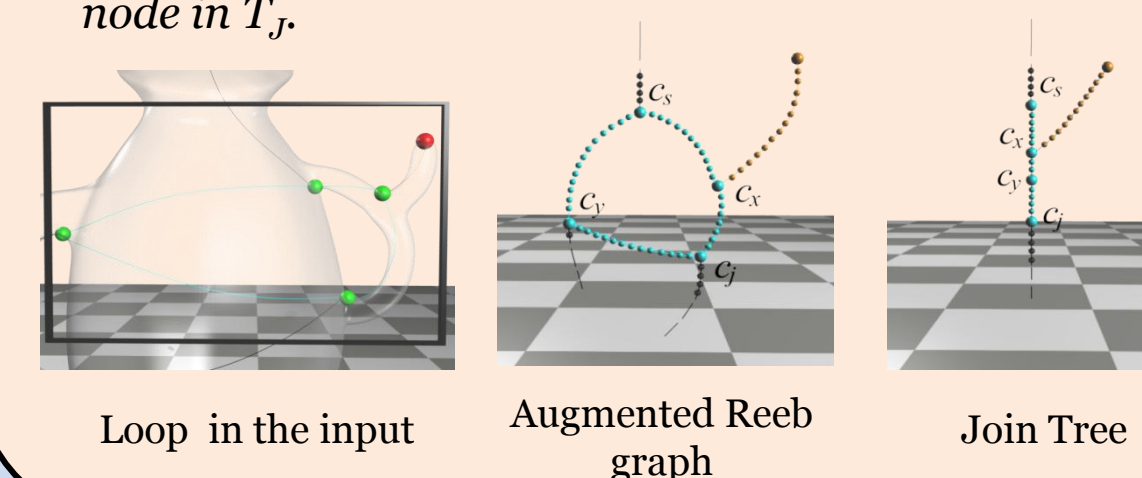
The ReCon Algorithm

1. Identify the loop saddles of the input
2. Split the input at a function value infinitesimally above that of the loop saddles to obtain a set of simply connected interval volumes
3. Compute the contour trees for each interval volume
4. Construct the Reeb graph by computing the union of these contour trees

Identifying Loop Saddles

Lemma

Let G_R be the Reeb graph of the given input scalar function f . Consider the join tree T_j of f . Any join saddle that ends a loop in G_R appears as a degree-2 node in T_j .



Properties

1. Efficient: has a running time of $O(n \log n + sn)$
 - n - # triangles
 - s - # saddles
2. Generic: works without any modifications on d -manifolds and non-manifolds
3. Easy to implement
4. Handles data that do not fit in memory
5. At least an order of magnitude faster than existing generic algorithms

Performance for 3D Input

3D Model	# Triangles	# Loops	Time taken (sec)			
			ReCon	LS	OS	Rand
Skull	0.34M	2	0.1	0.3	3.4	2.2
Post	1.24M	0	0.4	0.9	13.0	14.5
Plasma	2.64M	0	1.5	1.9	396.3	132.6
SF Earthquake	4.19M	0	2.4	2.8	598.1	166.9

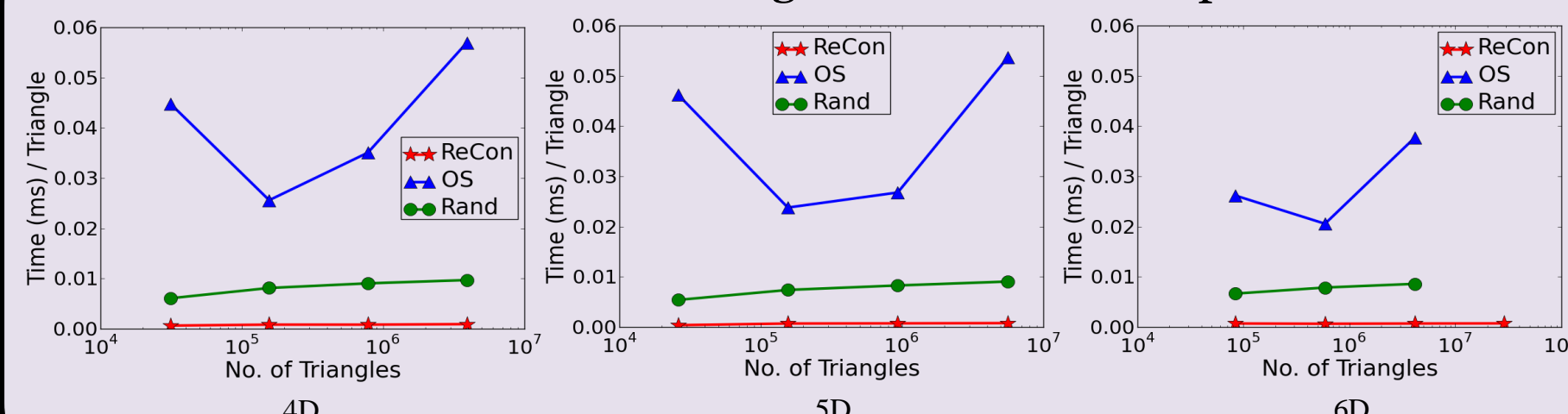
Performance for 2D Input

2D Model	# Triangles	# Loops	Time taken (sec)			
			ReCon	Online	OS	Rand
Awakening	4M	1643	2.4	6.72	41.4	51.8
Day	6M	2161	4.3	10.3	91.3	67.5
Dawn	6.6M	757	4.5	11.5	73.2	71.8
Lucy	28M	15	34.2	60.1	Mem	Mem

Scalability with increasing loop size

Model	# Triangles	# Loops	Time taken (sec)
s4d-7	0.78M	1×10^5	0.7
s5d-6	0.93M	1×10^5	0.7
s4d-8	3.9M	5.8×10^5	3.6
s5d-7	5.6M	5.6×10^5	4.4
s6d-6	4.1M	2.9×10^5	2.9
Lucy	28.0M	15	34.2
s6d-7	28.8M	2×10^6	21.2

Performance for higher dimensional input



Input – 4D, 5D, and 6D Sierpinski simplexes of various sizes

Handling Large data

Model	# Triangles	Function	Time taken	
			ReCon	Online
David	56M	x	3.6 m	4.7 m
		y	3.8 m	4.8 m
		z	3.2 m	16.6 m
St. Matthew	372M	x	26.9 m	40 m
		y	26.7 m	4.2 hrs
		z	25.2 m	41 m
Atlas	507M	x	41.5 m	*
		y	38.5 m	*
		z	42.6 m	*

References

- Pascucci et al. *ACM Trans. Graph* 2007 (Online)
- Tierny et al. *IEEE TVCG* 2009 (LS)
- Doraiswamy et al. *IEEE TVCG* 2011 (OS)
- Harvey et al. *SCG* 2010 (Rand)

Contact

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Acknowledgements

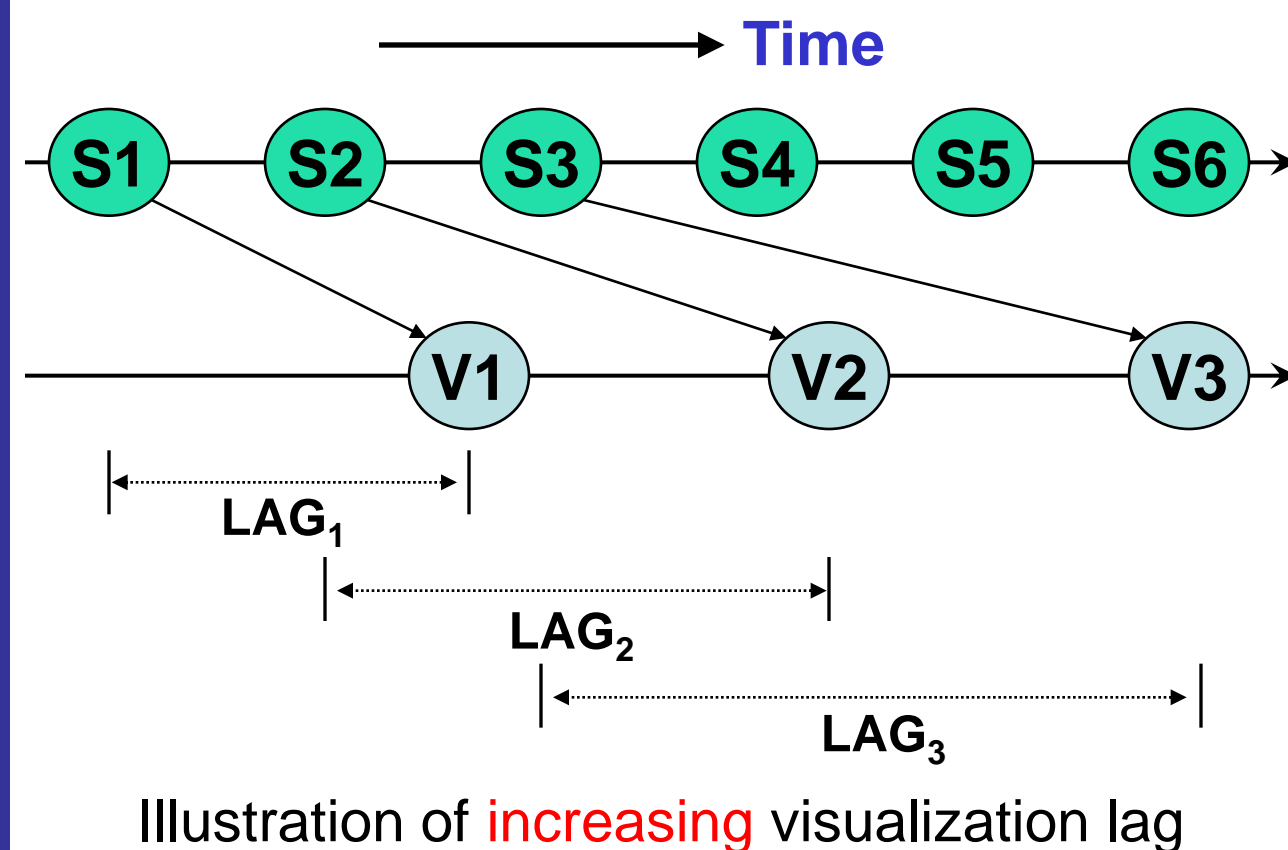
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Problem Statement

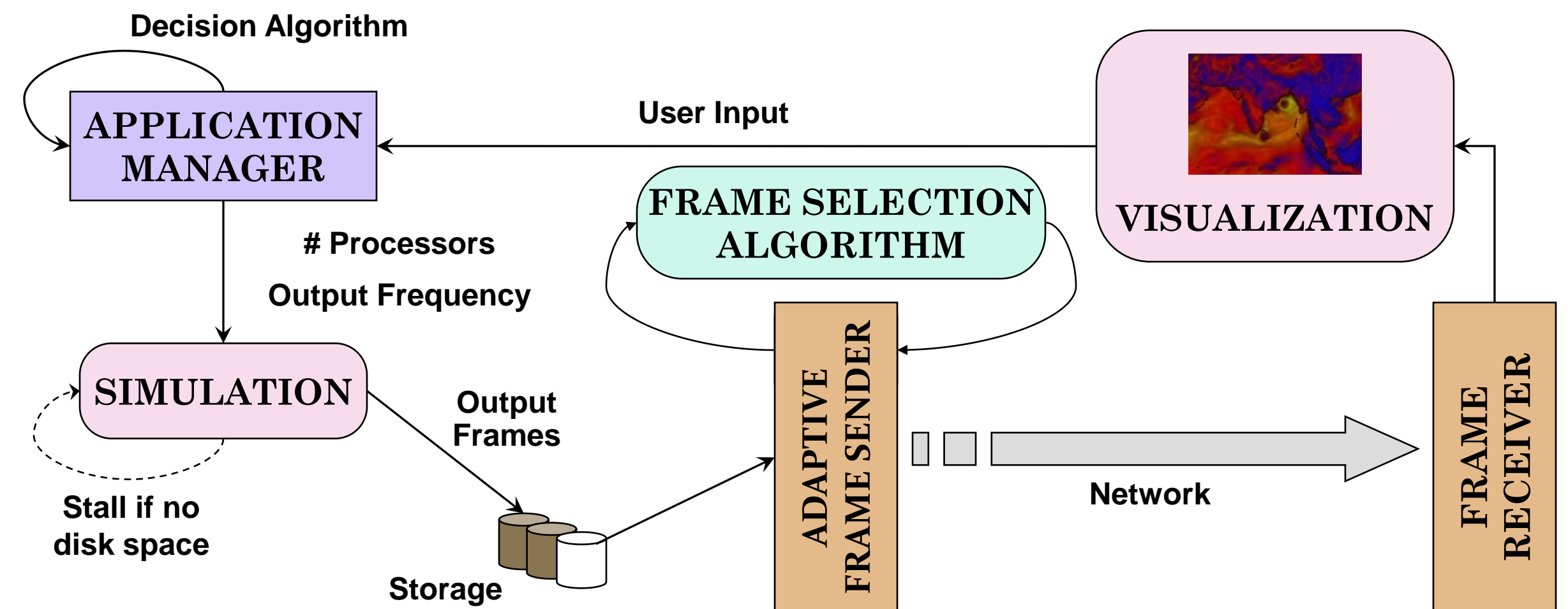
Develop an adaptive integrated steering framework that

- allows simultaneous simulation and online visualization
- spawns high-resolution simulation dynamically over desired region-of-interest
- supports optimal processor allocation for simulation
- supports optimal frequency of output for visualization

Simulation-Visualization Lag



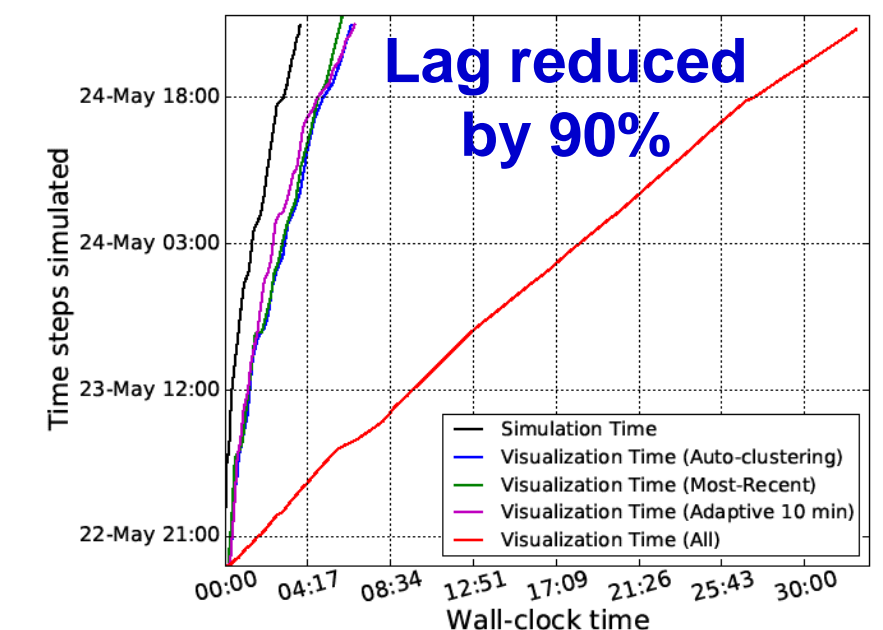
Framework for Adaptive Simulation and Online Visualization



Frame Selection Algorithms

- Most-recent
 - Transfer the most-recently simulated frame
- Auto-clustering
 - Modified k-means for temporal clustering to select the most-representative frames
- Adaptive
 - Transfer full or reduced frames within acceptable lag bound

Results

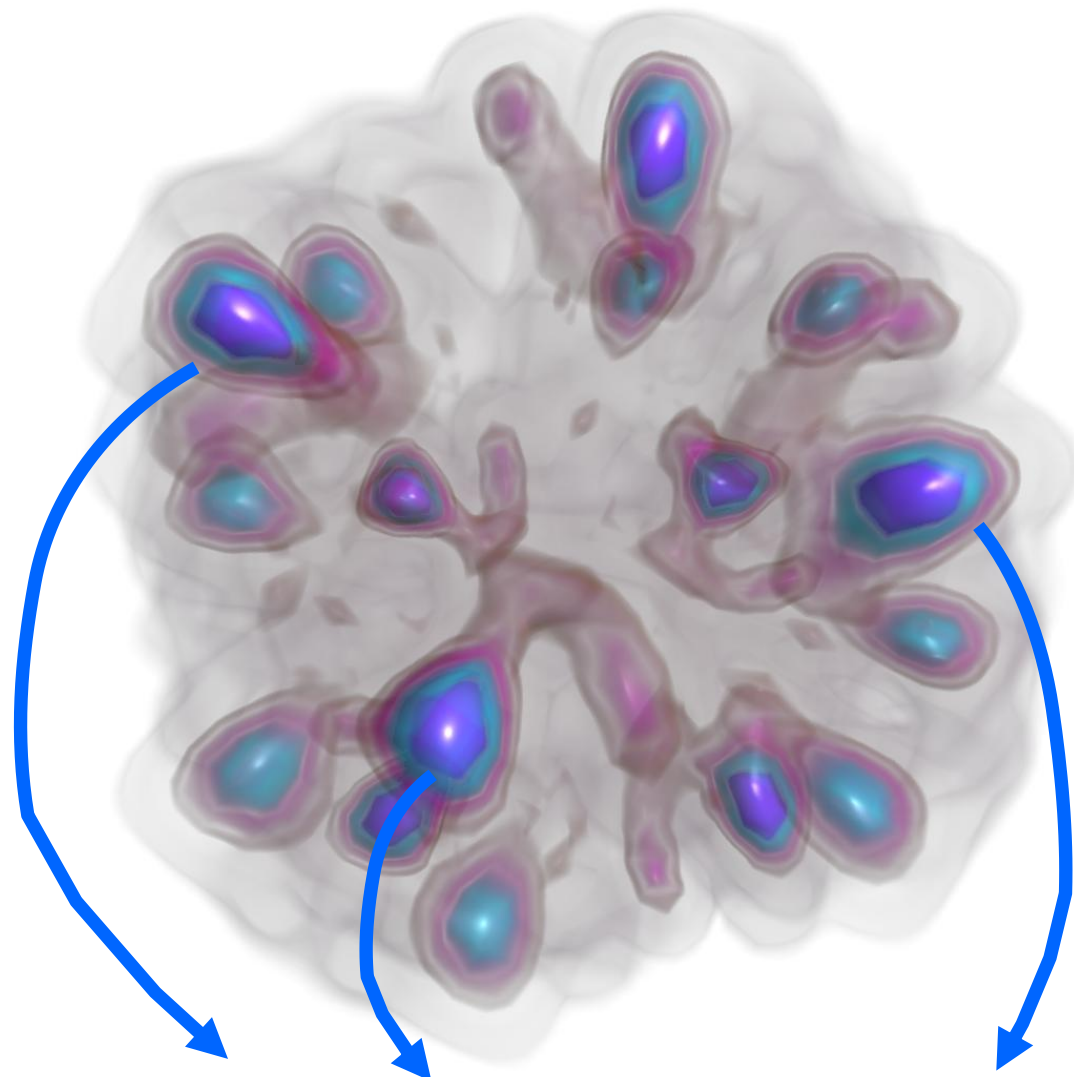




Symmetry in Scalar Field Topology

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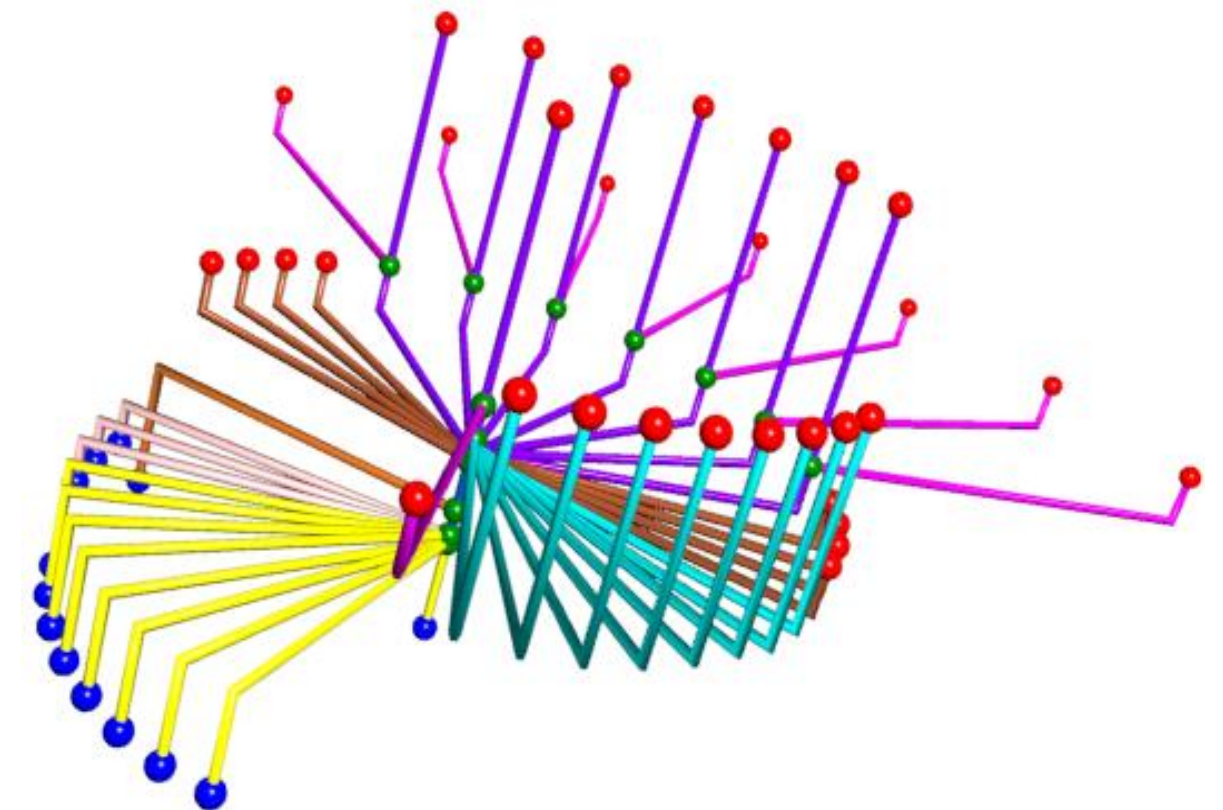
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Symmetric structures in EM image
of RuBisCO molecule



Four types of symmetric
regions identified



Classify subtrees of
the contour tree

Motivation

Scalar fields contain repeating patterns
Provide insights on scientific phenomena

Technique

Classify subtrees of the contour tree
Similarity measure to compare subtrees

Applications

Symmetry-aware isosurface extraction
Symmetry-aware transfer function design