## E0 234 : Homework 5

Deadline : 7th April 2023, 2pm

## Instructions

- Please write your answers using LATEX. Handwritten answers will not be accepted.
- You are forbidden from consulting the internet. You are strongly encouraged to work on the problems on your own.
- You may discuss these problems with another student. However, you must write your own solutions and must mention your collaborator's name. Otherwise, it will be considered as plagiarism.
- Academic dishonesty/plagiarism will be dealt with severe punishment.
- Late submissions are accepted only with prior approval (on or before the day of posting of HW) or a medical certificate.

## 1. Matrix rank in parallel.

Assume the following.

Given an  $n \times n$  integer matrix A with the property that rank  $A^2 = \operatorname{rank} A$  (the rank is computed over rational numbers), one can determine rank A using an efficient parallel algorithm.

We would like to use the above parallel algorithm to obtain an efficient randomized algorithm for computing matrix rank in general.

- (a) Pick an  $n \times n$  matrix R at random (by picking each entry of R independently from  $\{1, \ldots, n^2\}$ .
- (b) Determine the rank of RA using the above algorithm.

To establish that this algorithm determines the rank of A correctly with high probability proceed as follows.

- (a) Show that whp R has full rank, so ker RA = ker A.
- (b) Show that whp ker RA and im RA have only the vector **0** in common. (Hint: Let B consists of rankA linearly independent columns of A, and let C be a matrix whose columns form a basis for ker A. Consider det [RB | C] as a polynomial in the entries of R.)
- (c) Conclude that rank  $A = \operatorname{rank} RA = \operatorname{rank} (RA)^2$ .
- 2. Pattern matching.

Consider the following matrices:  $M(0) = \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix}$ ,  $M(1) = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}$ . For a string  $x \in \{0, 1\}^n$ , let  $M(x) = M(x_1)M(x_2)\cdots M(x_n)$ , where  $x = x_1x_2\ldots x_n$ . Show that this transformation has the following properties:

•  $M(x) = M(y) \Rightarrow x = y;$ 

• For  $x \in \{0,1\}^n$ , the entries of M(x) are bounded by the Fibonacci number  $F_n$ .

By considering the matrices M(x) modulo a randomly chosen prime in a suitable range, show how you would perform efficient randomized pattern matching. You may assume that the number of primes smaller than  $\tau$  is  $\pi(\tau) \sim t/\ln \tau$ .

3. Equivalent branching programs.

A (counting) branching program P on  $\mathbf{X} = \{X_1, X_2, \ldots, X_n\}$  is a directed acyclic graph with one source s and one sink t, where each edge is labelled by either a variable or its negation and no literal appears twice on any s-t path. The value P(x) of the program on an input  $x \in \{0, 1\}^n$  is the number of paths from s to t for which all literals appearing on the edges of the path evaluate to 1 under the assignment x. Describe a randomized algorithm that given two such branching programs P and Qefficiently determines (with small error) if P(x) = Q(x) for all  $x \in \{0, 1\}^n$ .

4. VC-Dimension.

Find the VC dimensions of the following range spaces. (a)  $S' := (X, \mathcal{R}')$ , where  $X = \mathbb{R}^2$  and  $\mathcal{R}'$  is the set of triangles. (b)  $S' := (X, \mathcal{R}')$ , where  $X = \mathbb{R}^2$  and  $\mathcal{R}'$  is the set of convex polygons with k sides.

5.  $\epsilon$ -net.

Let  $S := (X, \mathcal{R})$  be a range space with VC-dimension d. Prove that if a random sample M gives an  $\epsilon$ -net for S with probability at least  $1 - \delta$ , then |M| is  $\Omega(d/\epsilon)$ .

## Recommended practice problems (not for grading)

1. Book: Mitzenmacher-Upfal (2nd edition): 14.1, 14.2, 14.3, 14.4, 14.5, 14.7, 14.10, 14.11, 14.12.