

E0 224: Computational Complexity Theory
Indian Institute of Science
Assignment 3

Due date: Jan 15, 2021

Total marks: 50

1. **(12 marks)** Prove the following:
 - (a) **(3 marks)** Class PP has a complete problem under polynomial time Karp reduction.
 - (b) **(4 marks)** If $\text{BPP} = \text{PP}$ then PH collapses.
 - (c) **(5 marks)** PP is closed under complementation.
2. **(7 marks)** Prove that BP.NP is in Σ_3 .
3. **(9 marks)** Prove that $\overline{\text{SAT}} \in \text{BP.NP}$ implies $\text{PH} = \Sigma_3$.
4. **(4 marks)** Give a polynomial time algorithm that checks whether a given bipartite graph $G = (V, E)$ is contained in $\oplus \text{Perfect Matchings}$, where $\oplus \text{Perfect Matchings}$ is the set of all bipartite graphs having odd number of perfect matchings.
5. **(6 marks)** Prove that for any $n \times n$ matrix $A = (a_{i,j})_{i,j \in [n]}$,

$$\text{perm}(A) = \sum_{S \subseteq [n]} (-1)^{n-|S|} \prod_{i \in [n]} \left(\sum_{j \in S} a_{i,j} \right).$$

Use this to design an algorithm to compute the permanent in time $2^n \cdot \text{poly}(n)$.

6. **(4 marks)** Consider the following problem: Given an n -variate polynomial f in the form $\prod_{i \in [n]} \sum_{j \in [n]} a_{i,j} x_j$, where $a_{i,j}$ are integers, and $e_1, \dots, e_n \in \mathbb{Z}_{\geq 0}$ s.t. $e_1 + \dots + e_n = n$, compute

$$\frac{\partial^n f}{\partial x_1^{e_1} \partial x_2^{e_2} \dots \partial x_n^{e_n}}.$$

Prove that the problem is #P-hard.

7. **(8 marks)** Prove that #2-SAT is #P-complete.