## E0 224: Computational Complexity Theory Indian Institute of Science Assignment 2

## Due date: Oct 22, 2021

Total marks: 50

- 1. (3 marks) Define polyL to be  $\cup_{c>0}$ SPACE(log<sup>c</sup> n). Steve's Class SC is defined to be the set of languages that can be decided by deterministic machines that run in polynomial time and log<sup>c</sup> n space for some c > 0. It is an open problem whether PATH  $\in$  SC. Why does Savitch's theorem not resolve this question? Is SC the same as polyL  $\cap$  P.
- 2. (7 marks) Prove that in the read-once certificate definition of NL, if we allow the verifier machine to move its head back and forth on the certificate then the class being defined changes to NP.
- 3. (6 marks) If  $S = \{S_1, S_2, ..., S_m\}$  is a collection of subsets of a finite set U, the VC dimension of S, denoted VC(S), is the size of the largest set  $X \subseteq U$  such that for every  $X' \subseteq X$ , there is an *i* for which  $S_i \cap X = X'$ . (We say that X is *shattered* by S.)

A Boolean circuit C succinctly represents collections S if  $S_i$  consists of exactly those elements  $x \in U$  for which C(i, x) = 1. Finally,

VC-DIMENSION = {  $\langle C, k \rangle$  : C represents a collection S such that  $VC(S) \ge k$  }.

Show that VC-DIMENSION  $\in \Sigma_3$ .

- 4. (9 marks) Prove that a language L is in  $NC^1$  if and only if L is decided by a q(n)-size circuit family  $\{C_n\}_{n\in\mathbb{N}}$ , where q is a polynomial function and  $C_n$  is a Boolean formula for every  $n\in\mathbb{N}$ .
- 5. (10 marks) Linear programming (LP) is the problem of checking the feasibility of a system of linear inequality constraints over rationals. Prove that every language in P is logspace-reducible to LP. (In other words, LP is P-complete, and so, if LP is in NC, then P = NC.)
- 6. (6+9 marks) Prove that logspace uniform  $NC^1$  is contained in L. Prove that  $NL \subseteq NC$ .