## E0 224: Computational Complexity Theory Indian Institute of Science Assignment 2

Due date: Oct 20, 2023 Total marks: 50

- 1. (3 marks) Define polyL to be  $\cup_{c>0} \mathsf{SPACE}(\log^c n)$ . Steve's Class SC is defined to be the set of languages that can be decided by deterministic machines that run in polynomial time and  $\log^c n$  space for some c>0. It is an open problem whether PATH  $\in \mathsf{SC}$ . Why does Savitch's theorem not resolve this question? Is SC the same as  $\mathsf{polyL} \cap \mathsf{P}$ .
- 2. (7 marks) Prove that in the read-once certificate definition of NL, if we allow the verifier machine to move its head back and forth on the certificate then the class being defined changes to NP.
- 3. (6 marks) If  $S = \{S_1, S_2, ..., S_m\}$  is a collection of subsets of a finite set U, the VC dimension of S, denoted VC(S), is the size of the largest set  $X \subseteq U$  such that for every  $X' \subseteq X$ , there is an i for which  $S_i \cap X = X'$ . (We say that X is shattered by S.)

A Boolean circuit C succinctly represents collections S if  $S_i$  consists of exactly those elements  $x \in U$  for which C(i, x) = 1. Finally,

VC-DIMENSION =  $\{ \langle C, k \rangle : C \text{ represents a collection } S \text{ such that } VC(S) \geq k \}.$ 

Show that VC-DIMENSION  $\in \Sigma_3$ .

- 4. (9 marks) Prove that a language L is in  $NC^1$  if and only if L is decided by a q(n)-size circuit family  $\{C_n\}_{n\in\mathbb{N}}$ , where q is a polynomial function and  $C_n$  is a Boolean formula for every  $n\in\mathbb{N}$ .
- 5. (10 marks) Linear programming (LP) is the problem of checking the feasibility of a system of linear inequality constraints over rationals. Prove that every language in P is logspace-reducible to LP. (In other words, LP is P-complete, and so, if LP is in P-complete, and so, if LP is in P-complete, and so if P-complete, and P-complete, an
- 6. (6+9 marks) Prove that logspace uniform  $NC^1$  is contained in L. Prove that  $NL \subseteq NC$ .