#### Lecture 7: Singular Value Decomposition - II

November 8, 2017

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- Power Method, Fundamental Thm of Markov Chains

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- $\sum_{i=1}^{k} A \mathbf{v_i} \mathbf{v_i}^T = \sum_{i=1}^{k} \sigma_i \mathbf{u_i} \mathbf{v_i}^T = A_k$ , proving Lemma.

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- Will also prove  $A_k$  is best rank k approx in spectral norm.



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- For  $\varepsilon>0$ , let  $\mathbf{v}_{\mathbf{i}}'=\frac{\mathbf{v_i}+\epsilon\mathbf{v_j}}{|\mathbf{v_i}+\epsilon\mathbf{v_j}|}$ .  $|\mathbf{v}_{\mathbf{i}}'|=1$  and  $A\mathbf{v}_{\mathbf{i}}'=\frac{\sigma_i\mathbf{u_i}+\varepsilon\sigma_j\mathbf{u_j}}{\sqrt{1+\varepsilon^2}}$  has length at least as large as its component along  $\mathbf{u_i}$  which is

$$\mathbf{u_i^T}\left(\frac{\sigma_i\mathbf{u_1}+\varepsilon\sigma_i\mathbf{u_i}}{\sqrt{1+\varepsilon^2}}\right) > \left(\sigma_i+\varepsilon\sigma_i\delta\right)\left(1-\frac{\varepsilon^2}{2}\right) > \sigma_i-\frac{\varepsilon^2}{2}\sigma_i+\varepsilon\sigma_i\delta-\frac{\varepsilon^3}{2}\sigma_i\delta > \sigma_i, \text{ for sufficiently small } \epsilon, \text{ a contradiction since } \mathbf{v_i}+\varepsilon\mathbf{v_j} \text{ is orthogonal to } \mathbf{v_1},\mathbf{v_2},\ldots,\mathbf{v_{i-1}} \text{ since } j>i \text{ and } \sigma_i \text{ is defined to be the maximum of } |A\mathbf{v}| \text{ over such vectors.}$$

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- $\bullet |(A A_k)\mathbf{v}| = \left| \sum_{i=k+1}^r \sigma_i \mathbf{u}_i \mathbf{v}_i^T \sum_{j=1}^r c_j \mathbf{v}_j \right| = \left| \sum_{i=k+1}^r c_i \sigma_i \mathbf{u}_i \mathbf{v}_i^T \mathbf{v}_i \right| = \left| \sum_{i=k+1}^r c_i \sigma_i \mathbf{u}_i \right| = \sqrt{\sum_{i=k+1}^r c_i^2 \sigma_i^2}.$

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- For **v** maximizing this, subject to  $|\mathbf{v}|^2 = \sum_{i=1}^r c_i^2 = 1$ , have  $c_{k+1} = 1$ , rest  $c_i = 0$ . QED



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- Since  $z \in \text{Span}\{v_1, v_2, \dots, v_{k+1}\}$

$$|\mathbf{A}\mathbf{z}|^2 = \left|\sum_{i=1}^n \sigma_i \mathbf{u_i} \mathbf{v_i}^T \mathbf{z}\right|^2 = \sum_{i=1}^n \sigma_i^2 \left(\mathbf{v_i}^T \mathbf{z}\right)^2 = \sum_{i=1}^{k+1} \sigma_i^2 \left(\mathbf{v_i}^T \mathbf{z}\right)^2 \ge$$

$$\sigma_{k+1}^2 \sum_{i=1}^{k+1} \left( \mathbf{v_i}^T \mathbf{z} \right)^2 = \sigma_{k+1}^2$$
. QED

