

# Gödel's Incompleteness Theorem

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# Outline

- 1 Theory of Arithmetic
- 2 Peano's Proof System for Arithmetic
- 3 Proof of Gödel's theorem

# Gödel's Incompleteness Theorem

## Theorem (Gödel)

*There **cannot** exist a sound and complete proof system for arithmetic (i.e. Natural numbers with addition and multiplication  $(\mathbb{N}, +, \cdot)$ )*

# Arithmetic

First-order logic of  $(\mathbb{N}, +, \cdot)$ :

- Domain is  $\mathbb{N} = \{0, 1, 2, \dots\}$ .
- Terms:  $0$ ,  $1$ ,  $0 + 1$ ,  $1 \cdot x$ ,  $x + y$ ,  $x \cdot y$ , etc.
- Atomic formulas:  $t = t$
- Note that relations like “ $<$ ” are definable in the logic:  $t < t'$  is definable as  $\exists x(t + x = t')$ .
- Formulas:
  - Atomic formulas
  - Quantification:  $\forall x\varphi$ ,  $\exists x\varphi$
  - Boolean combinations:  $\neg\varphi$ ,  $\varphi \vee \psi$ ,  $\varphi \wedge \psi$ .

# What we can say in $\text{FO}(\mathbb{N}, +, \cdot)$

- “Integer division of  $x$  by  $y$  gives quotient  $q$  and leaves remainder  $r$ ”

$$\text{intdiv}(x, y, q, r) \stackrel{\text{def}}{=} x = (q \cdot y) + r \wedge r < y.$$

- “ $y$  divides  $x$ ”

$$\text{divides}(y, x) \stackrel{\text{def}}{=} \exists q (x = q \cdot y).$$

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- “ $x$  is a power of 2”

$$\text{power}_2(x) \stackrel{\text{def}}{=} \forall p ((\text{prime}(p) \wedge \text{divides}(p, x)) \implies p = 2).$$

# What we can say in $\text{FO}(\mathbb{N}, +, \cdot)$

- “Every number has a successor”

$$\forall n \exists m (m = n + 1).$$

- “Every number has a predecessor”

$$\forall n \exists m (n = m + 1).$$

- “There are only finitely many primes”

$$\exists n \forall p (\text{prime}(p) \implies p < n).$$

- “There are infinitely many primes”

$$\forall n \exists p (\text{prime}(p) \wedge p > n).$$

# Theory of $\text{FO}(\mathbb{N}, +, \cdot)$

$Th(\mathbb{N}, +, \cdot)$  is the set of sentences of  $\text{FO}(\mathbb{N}, +, \cdot)$  that are true. For example:

- “Every number has a successor”

$$\forall n \exists m (m = n + 1).$$

belongs to  $Th(\mathbb{N}, +, \cdot)$ , while

- “There are only finitely many primes”

$$\exists n \forall p (\text{prime}(p) \implies p < n).$$

does not.

Note that there is a mathematical definition of truth based on the mathematical definition of the semantics of the logic.

# Peano's Proof System for Arithmetic

- Axioms:

$$\begin{aligned}
 &\forall x \neg(0 = x + 1) \\
 &\forall x \forall y (x + 1 = y + 1 \implies x = y) \\
 &\forall x (x + 0 = x) \\
 &\forall x \forall y \forall z (x + (y + z) = (x + y) + z) \\
 &\forall x (x \cdot 0 = 0) \\
 &\forall x \forall y \forall z (x \cdot (y + z) = (x \cdot y) + (x \cdot z)) \\
 &(\varphi(0) \wedge \forall x (\varphi(x) \implies \varphi(x + 1))) \implies \forall x \varphi(x).
 \end{aligned}$$

- Other axioms like  $(\varphi \wedge \psi) \implies \varphi$ ,  $\forall x(\varphi) \implies \varphi(17)$ .
- Inference rules like 'Modus Ponens'

Given  $\varphi$  and  $\varphi \implies \psi$ , infer  $\psi$ .

# Proof

A proof of  $\varphi$  in a proof system is a finite sequence of sentences

$$\varphi_0, \varphi_1, \dots, \varphi_n$$

such that each  $\varphi_i$  is either an axiom or follows from two previous ones by an inference rule, and  $\varphi_n = \varphi$ .

A proof system is “sound” if whatever it proves is indeed true (i.e. in  $Th(\mathbb{N})$ ).

A proof system is “complete” if whatever it can prove whatever is true (i.e. in  $Th(\mathbb{N})$ ).

# Gödel's Incompleteness Theorem

## Theorem (Gödel)

*There **cannot** exist a sound and complete proof system for arithmetic (i.e. Natural numbers with addition and multiplication  $(\mathbb{N}, +, \cdot)$ )*

# Proof of Gödel's theorem

- Gödel's original proof was an intricate construction of an  $FO(\mathbb{N}, +, \cdot)$  sentence  $\varphi$  which (for a given proof system like Peano's) asserts that

*"I am not provable in the given proof system"*

- It follows that the proof system is either unsound (if  $\vdash \varphi$ ) or incomplete (if  $\nvdash \varphi$ ).
- Here we will follow a subsequent proof given by Turing which shows

$$\neg \text{HP} \leq Th(\mathbb{N}).$$

- Hence  $Th(\mathbb{N})$  is not even r.e. and hence there cannot be a proof system that is sound and complete (why?).

# Encoding computations of $M$ on $x$

Let  $M = (Q, A, \Gamma, s, \delta, \vdash, \flat, t, r)$  be a given TM and let  $x = a_1 a_2 \cdots a_n$  be an input to it.

We can represent a configuration of  $M$  as follows:

$$\begin{array}{ccccccc} \vdash & b_1 & b_2 & b_3 & \cdots & b_m \\ - & - & q & - & & - \end{array}$$

Thus a configuration is encoded over the alphabet  $\Gamma \times (Q \cup \{-\})$ .



# Encoding computations of $M$ on $x$

A computation of  $M$  on  $x$  is string of the form

$$c_0 \# c_1 \# \cdots \# c_N \#$$

such that

- ① Each  $c_i$  is the encoding of a configuration of  $M$ .
- ②  $c_0$  is (encoding of) the start configuration of  $M$  on  $x$ .

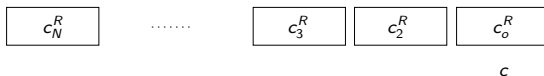
$$\begin{array}{ccccccc} \vdash & a_1 & a_2 & a_3 & \cdots & a_n \\ s & - & - & - & & - \end{array}$$

- ③ Each  $c_i \xRightarrow{1} c_{i+1}$ , and
- ④  $c_N$  is a halting configuration (i.e. state component is  $t$  or  $r$ ).



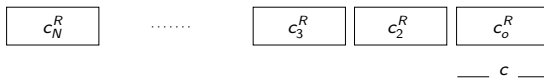
# Basic idea

View a computation of  $M$  on  $x$  as a number whose representation in base  $p \geq |\Delta|$  looks like:



Now construct a sentence  $\varphi_{M,x}$  which asserts that “there is a number  $n$  whose base- $p$  representation encodes a valid halting computation of  $M$  on  $x$ .”

# The sentence $\varphi_{M,x}$



- Define  $valcomp_{M,x}(v)$  to be

$$\begin{aligned} \exists c \exists d ( & power_p(c) \wedge power_p(d) \\ & \wedge length(v, d) \wedge start(v, c) \\ & \wedge move(v, c, d) \wedge halt(v, d)). \end{aligned}$$

- Define  $\varphi_{M,x}$  to be

$$\exists v \, valcomp_{M,x}(v).$$

# Expressing the components of $\varphi_{M,x}$

The key predicate we need is “ $digit_p(v, d, a)$ ”: which says that  $d$  is a power of  $p$  (say  $d = p^k$ ), and in the base- $p$  representation, the  $k$ -th digit is  $a$ .

$$digit_p(v, d, a) \stackrel{\text{def}}{=} \exists u \exists r (v = u \cdot p \cdot d + a \cdot d + r \wedge r < d).$$