More undecidable problems

Deepak D'Souza

Department of Computer Science and Automation Indian Institute of Science, Bangalore.

18 November 2019

▲ロト ▲帰ト ▲ヨト ▲ヨト 三日 - の々ぐ

Outline





Problem (a)

Is it decidable whether a given Turing machine has at least 481 states? Assume that the TM is given using the encoding below:

 $0''10''10^{k}10^{s}10^{t}10'10''10''10''10^{p}10^{a}10''10^{b}10110''10^{a'}10^{a'}10^{a'}10^{b'}100\cdots 10^{p''}10^{a''}10^{q''}10^{b''}10.$

Problem (a)

Is it decidable whether a given Turing machine has at least 481 states? Assume that the TM is given using the encoding below:

 $0''10''10^{k}10^{s}10^{t}10'10''10''10''10^{p}10^{a}10''10^{b}10110''10^{a'}10^{a'}10^{a'}10^{b'}100\cdots 10^{p''}10^{a''}10^{q''}10^{b''}10.$

Yes, it is.

We can give a TM N which given enc(M)

- Counts the number of states in *M* upto 481.
- Accepts if it reaches 481, rejects otherwise.

Problem (b)

Is it decidable whether a given Turing machine takes more than 481 steps on input ϵ without halting?

Problem (b)

Is it decidable whether a given Turing machine takes more than 481 steps on input ϵ without halting?

Yes, it is.

We can give a TM N which given enc(M)

- Uses 4 tapes: On the 4th tape it writes 481 0's.
- Uses the first 3 tapes to simulate M on input ϵ , like the universal TM U.
- Blanks out a 0 from 4th tape for each 1-step simulation done by *U*.
- Rejects if *M* halts before all 0's are blanked out on 4th tape, accepts otherwise.

Problem (c)

Is it decidable whether a given Turing machine takes more than 481 steps on *some* input without halting?

◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 のへぐ

Problem (c)

Is it decidable whether a given Turing machine takes more than 481 steps on *some* input without halting?

Yes, it is.

Check if M runs for more than 481 steps on some input x of length upto 481. If so accept, else reject.

Problem (d)

Is it decidable whether a given Turing machine takes more than 481 steps on *all* inputs without halting?

Problem (d)

Is it decidable whether a given Turing machine takes more than 481 steps on *all* inputs without halting?

Yes, it is.

Check if M runs for more than 481 steps on each input x of length upto 481. If so accept, else reject.

Problem (e)

Is it decidable whether a given Turing machine moves its head more than 481 cells away from the left-end marker, on input ϵ ?

◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 のへぐ

Problem (e)

Is it decidable whether a given Turing machine moves its head more than 481 cells away from the left-end marker, on input ϵ ?

Yes, it is. Simulate *M* on ϵ for upto $m^{481} \cdot 482 \cdot k + 1$ steps. If *M* visits the 482nd cell, accept, else reject.

Problem (f)

Is it decidable whether a given Turing machine accepts the null-string ϵ ?

Problem (f)

Is it decidable whether a given Turing machine accepts the null-string ϵ ?

No.

If it were decidable, say by a TM N, then we could use N to decide HP as follows: Define a new machine N' which given input M # x, outputs the description of a machine $M'_{M,x}$ which:

- erases its input
- writes x on its input tape
- Behaves like *M* on *x*
- Accepts if *M* halts on *x*.

N' then calls N with input $M'_{M,x}$. N accepts $M'_{M,x}$ iff $M'_{M,x}$ accepts ϵ iff M halts on x.

Turing machine M' for Problem (f)

$$L(M'_{M,x}) = \begin{cases} A^* & \text{if } M \text{ halts on } x \\ \emptyset & \text{if } M \text{ does not halt on } x. \end{cases}$$

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 のへぐ

More decidable/undecidable problems

Problem (g)

Is it decidable whether a given Turing machine accepts any string at all? That is, is $L(M) \neq \emptyset$?

More decidable/undecidable problems

Problem (h)

Is it decidable whether a given Turing machine accepts all strings? That is, is $L(M) = A^*$?

More decidable/undecidable problems

Problem (i)

Is it decidable whether a given Turing machine accepts a finite set?

◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 のへぐ

Problem (j)

Is it decidable whether a given Turing machine accepts a regular set?

Problem (j)

Is it decidable whether a given Turing machine accepts a regular set?

Given M and x, build a new machine M' that behaves as follows:

- Saves its input y on tape 2.
- writes x on tape 1.
- \bigcirc runs as M on x.
- if *M* gets into a halting state, then
 - M' takes back control,
 - Runs as M_{NR} on y,
 - (Here M_{NR} is any TM that accepts a non-regular language NR, say NR = {aⁿbⁿ | n ≥ 0}).
 - M' accepts iff M_{NR} accepts.

Turing machine M' for Problem (j)

$$L(M') = \begin{cases} NR & \text{if } M \text{ halts on } x \\ \emptyset & \text{if } M \text{ does not halt on } x. \end{cases}$$

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 のへぐ

More decidable/undecidable problems

Problem (k)

Is it decidable whether a given Turing machine accepts a CFL?

◆□▶ ◆□▶ ◆三▶ ◆三▶ 三三 のへぐ

More decidable/undecidable problems

Problem (I)

Is it decidable whether a given Turing machine accepts a recursive set?