Floyd-Hoare Style Program Verification
FMSE Course

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28 Feb 2022

Outline of these lectures

Checking Pre/Post Assertions in Programs

- Moving on from reasoning about models to reasoning about code.
- Still a deductive style of verification.
- Helps us to verify assertions and also refinement-based functionality verification.

Floyd-Hoare Style of Program Verification

Robert W. Floyd: "Assigning meanings to programs" *Proceedings of the American Mathematical Society Symposia on Applied Mathematics* (1967)

C A R Hoare: "An axiomatic basis for computer programming", *Communications of the ACM* (1969).

Floyd-Hoare Logic

- A way of asserting properties of programs.
- Hoare triple: $\{A\}P\{B\}$ asserts that "Whenever program P is started in a state satisfying condition A, if it terminates, it will terminate in a state satisfying condition B."
- Example assertion: $\{n \geq 0\}\ P\ \{a = n + m\}$, where P is the program:

```
int a := m;
int x := 0;
while (x < n) {
  a := a + 1;
  x := x + 1;
}
```

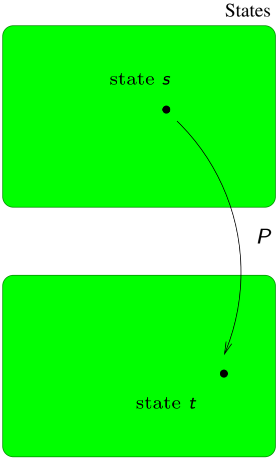
- Inductive Annotation ("consistent interpretation") (due to Floyd)
- A proof system (due to Hoare) for proving such assertions.
- A way of reasoning about such assertions using the notion of "Weakest Preconditions" (due to Dijkstra).

A simple programming language

- `skip` (do nothing)
- `x := `e (assignment)
- `if` b `then` S `else` T (if-then-else)
- `while` b `do` S (while loop)
- S ; T (sequencing)

Programs as State Transformers

- Program state is valuation to variables of the program: $States = Var \to \mathbb{Z}$.
- View program P as a partial map $[\![P]\!] : States \to States$.

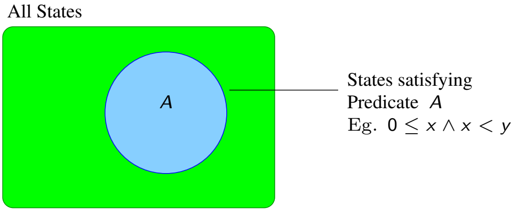


$s : \langle x \mapsto 2,\ y \mapsto 10,\ z \mapsto 3 \rangle$

```
y := y + 1;
z := x + y
```

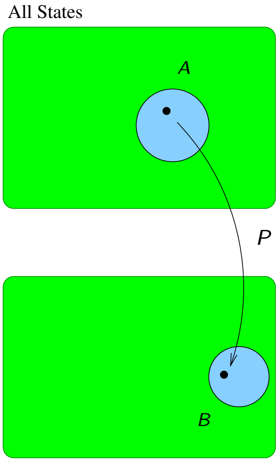
$t : \langle x \mapsto 2,\ y \mapsto 11,\ z \mapsto 13 \rangle$

Predicates on States



Assertion of "Partial Correctness" $\{A\}P\{B\}$

$\{A\}P\{B\}$ asserts that "Whenever program P is started in a state satisfying condition A, either it will not terminate, or it will terminate in a state satisfying condition B."



$\{10 \leq y\}$

```
y := y + 1;
z := x + y
```

$\{x < z\}$

Mathematical meaning of a Hoare triple

- View program P as a relation on States (allows non-termination as well as non-determinism)

$$[\![P]\!] \subseteq \mathrm{States} \times \mathrm{States}.$$

  Here $(s, t) \in [\![P]\!]$ iff it is possible to start P in the state s and terminate in state t.
- $[\![P]\!]$ is possibly non-determinisitic, in case we also want to model non-deterministic assignment etc.
- Then the Hoare triple $\{A\}\ P\ \{B\}$ is true iff for all states s and t: whenever $s \models A$ and $(s, t) \in [\![P]\!]$, then $t \models B$.
- In other words $Post_{[\![P]\!]}([\![A]\!]) \subseteq [\![B]\!]$.

Example programs and pre/post conditions

```
// Pre: true

if (a <= b)
  min := a;
else
  min := b;

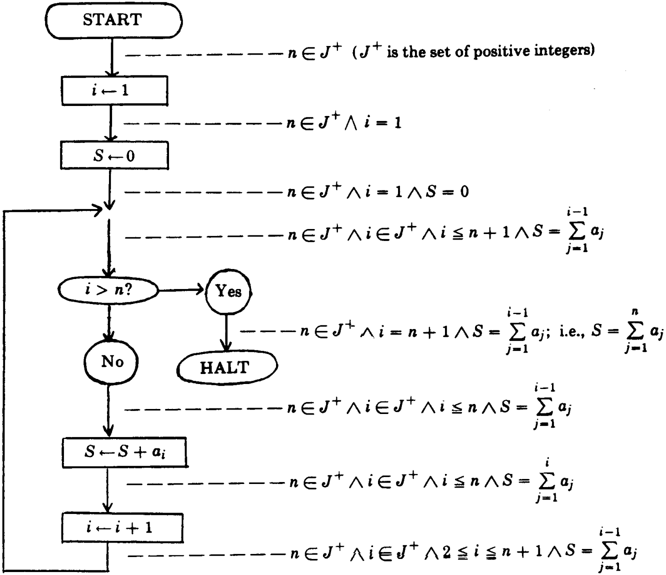
// Post: min <= a && min <= b
```

```
// Pre: 0 <= n

int a := m;
int x := 0;
while (x < n) {
  a := a + 1;
  x := x + 1;
}

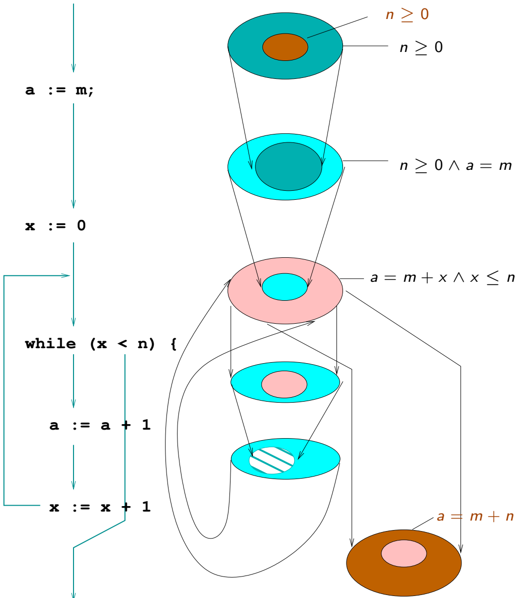
// Post: a = m + n
```

Floyd style proof: Inductive Annotation



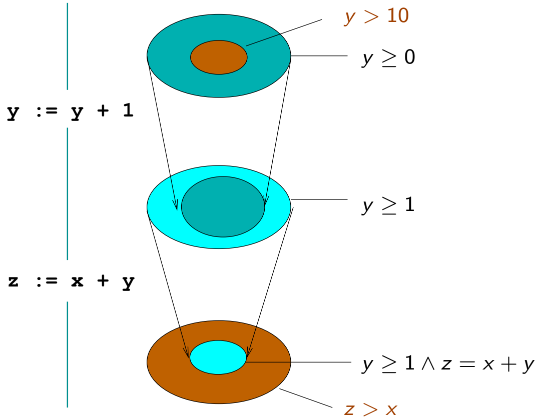
Inductive annotation based proof of a pre/post specification

- Annotate each program point i with a predicate A_i
- Successive annotations must be inductive: $[\![S_i]\!]([\![A_i]\!]) \subseteq [\![A_{i+1}]\!]$, OR logically: $A_i \wedge [S_i] \implies A'_{i+1}$.
- Annotation is adequate: $Pre \implies A_1$ and $A_n \implies Post$.
- Adequate annotation constitutes a proof of $\{Pre\}$ Prog $\{Post\}$.



Example of inductive annotation

To prove: $\{y > 10\}$ `y := y+1; z := x+y` $\{z > x\}$



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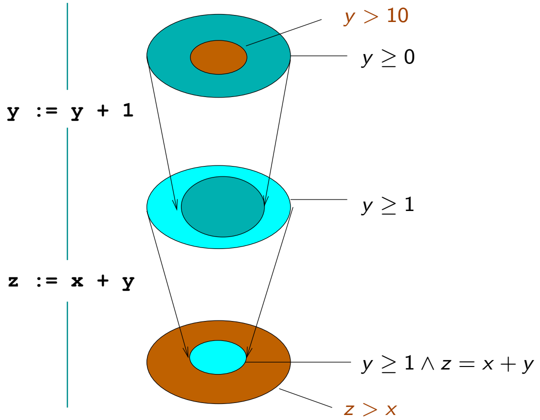


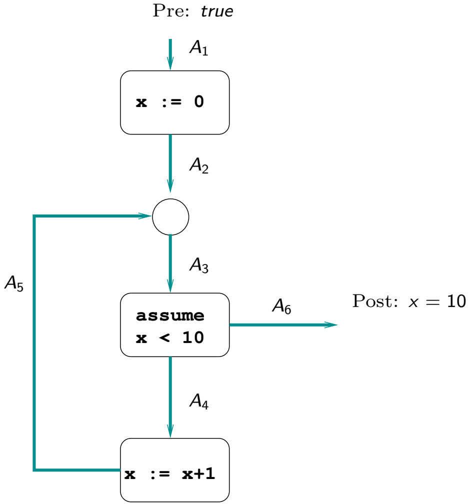
Logical proof obligations (VCs):

$(y > 10 \implies y \geq 0) \wedge ((y \geq 1 \wedge z = x + y) \implies z > x) \wedge$
$((y \geq 0 \wedge y' = y + 1) \implies y' \geq 1) \wedge ((y \geq 1 \wedge z' = x + y) \implies y' \geq$

Exercise 1

Prove using Floyd-style annotation:

```
// Pre: true
int x := 0;
while (x < 10)
  x := x + 1;
// Post: x = 10
```



Also write out the proof obligations (verification conditions).

Exercise 2

Prove using Floyd's inductive annotation:

$$\{n \geq 1\}\ P\ \{a = n!\},$$

where P is the program:

```
x := n;
a := 1;
while (x ≥ 1) {
    a := a * x;
    x := x - 1
}
```

Assume that factorial is defined as follows:

$$n! = \begin{cases} n \times (n-1) \times \cdots \times 1 & \text{if} \quad n \geq 1 \\ 1 & \text{if} \quad n = 0 \\ -1 & \text{if} \quad n < 0 \end{cases}$$

Exercise 2

Prove using Floyd's inductive annotation:

$$\{n \geq 1\}\ P\ \{a = n!\},$$

where P is the program:

```
S1:  x := n;
S2:  a := 1;
S3:  while (x ≥ 1) {
S4:      a := a * x;
S5:      x := x - 1
     }
```

Assume that factorial is defined as follows:

$$n! = \begin{cases} n \times (n-1) \times \cdots \times 1 & \text{if} \quad n \geq 1 \\ 1 & \text{if} \quad n = 0 \\ -1 & \text{if} \quad n < 0 \end{cases}$$

Hoare's view: Program as a composition of statements

```
int a := m;
int x := 0;
while (x < n) {
  a := a + 1;
  x := x + 1;
}
```

Hoare's view: Program as a composition of statements

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int a := m;
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  a := a + 1;
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}
```

```
S1: int a := m;
S2: int x := 0;
S3: while (x < n) {
      a := a + 1;
      x := x + 1;
    }
```

Program is `S1;S2;S3`

Proof rules of Hoare Logic

Axiom of Valid formulas:

$$\frac{}{A}$$

provided "$\models A$" (i.e. A is a valid logical formula, eg. $x > 10 \implies x > 0$).

Skip:

$$\frac{}{\{A\}\ \texttt{skip}\ \{A\}}$$

Assignment

$$\frac{}{\{A[e/x]\}\ \texttt{x := e}\ \{A\}}$$

Proof rules of Hoare Logic

If-then-else:

$$\frac{\{P \wedge b\}\ S\ \{Q\},\ \{P \wedge \neg b\}\ T\ \{Q\}}{\{P\}\ \texttt{if}\ \ b\ \ \texttt{then}\ \ S\ \ \texttt{else}\ \ T\ \{Q\}}$$

While (here P is called a *loop invariant*)

$$\frac{\{P \wedge b\}\ S\ \{P\}}{\{P\}\ \texttt{while}\ \ b\ \ \texttt{do}\ \ S\ \{P \wedge \neg b\}}$$

Sequencing:

$$\frac{\{P\}\ S\ \{Q\},\ \{Q\}\ T\ \{R\}}{\{P\}\ S\,;T\ \{R\}}$$

Weakening:

$$\frac{P \implies Q,\ \{Q\}\ S\ \{R\},\ R \implies T}{\{P\}\ S\ \{T\}}$$

Loop invariants

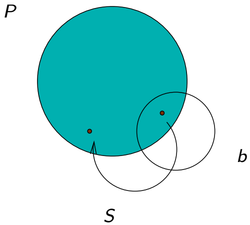
A predicate P is a loop invariant for the while loop:

```
while (b) {
  S
}
```

if $\{P \land b\}\ S\ \{P\}$ holds.

If P is a loop invariant then we can infer that:

$$\{P\}\ \texttt{while}\ b\ \texttt{do}\ S\ \{P \land \neg b\}$$



Some examples to work on

Use the rules of Hoare logic to prove the following assertions:

1. $\{x > 3\}$ `x := x + 2` $\{x \geq 5\}$
2. $\{(y \leq 0) \wedge (-1 < x)\}$ `if` $(y < 0)$ `then x:=x+1 else x:=y` $\{0 \leq x\}$
3. $\{x \leq 0\}$ while $(x \leq 5)$ do x := x+1 $\{x = 6\}$

Example proof using Hoare Logic

1. $\{n \geq 0\}$ S1 $\{n \geq 0 \wedge a = m\}$
2. $\{n \geq 0 \wedge a = m\}$ S2 $\{n \geq 0 \wedge a = m \wedge x = 0\}$
3. $\{a = m + x \wedge 0 \leq x \leq n \wedge x < n\}$ S4;S5 $\{a = m + x \wedge 0 \leq x \leq n\}$
4. $\{a = m + x \wedge 0 \leq x \leq n\}$ S3 $\{a = m + x \wedge 0 \leq x \leq n \wedge x \geq n\}$
5. $\{n \geq 0\}$ S1;S2 $\{n \geq 0 \wedge a = m \wedge x = 0\}$ (From Seq rule, 1 and 2)
6. $(n \geq 0 \wedge a = m \wedge x = 0) \implies (a = m + x \wedge 0 \leq x \leq n)$ (From logical axiom)
7. $\{n \geq 0\}$ S1;S2 $\{a = m + x \wedge 0 \leq x \leq n\}$ (From Weakening rule, 5 and 6)
8. $\{n \geq 0\}$ (S1;S2);S3 $\{a = m + x \wedge 0 \leq x \leq n \wedge x \geq n\}$ (From Seq rule, 7, 4)
9. $(a = m + x \wedge 0 \leq x \leq n \wedge x \geq n) \implies (a = m + n)$
10. $\{n \geq 0\}$ (S1;S2);S3 $\{a = m + n\}$ (From Weakening rule, 8, 9).

```
// pre: n >= 0
S1: int a := m;
S2: int x := 0;
S3: while (x < n) {
S4:   a := a + 1;
S5:   x := x + 1;
    }
// post: a = m + n
```

Program is `S1;S2;S3`

Exercise

Prove using Hoare logic:

$$\{n \geq 1\}\ P\ \{a = n!\},$$

where P is the program:

```
x := n;
a := 1;
while (x ≥ 1) {
   a := a * x;
   x := x - 1
}
```

Assume that factorial is defined as follows:

$$n! = \begin{cases} n \times (n-1) \times \cdots \times 1 & \text{if} \quad n \geq 1 \\ 1 & \text{if} \quad n = 0 \\ -1 & \text{if} \quad n < 0 \end{cases}$$

Exercise

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     }
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Assume that factorial is defined as follows:

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Soundness and Completeness

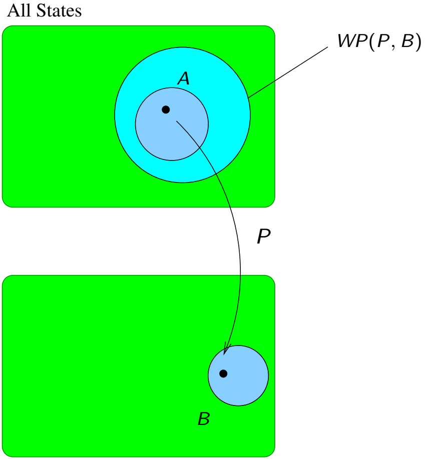
Soundness: If our proof system proves $\{A\}\ P\ \{B\}$ then $\{A\}\ P\ \{B\}$ indeed holds.

Completeness: If $\{A\}\ P\ \{B\}$ is true then our proof system can prove $\{A\}\ P\ \{B\}$.

- Floyd proof style is sound since any execution must stay within the annotations. Complete because the "collecting" set is an adequate inductive annotation for any program and any true pre/post condition.
- Hoare logic is sound, essentially because the individual rules can be seen to be sound.
- For completness of Hoare logic, we need weakest preconditions.

Weakest Precondition $WP(P, B)$

$WP(P, B)$ is "a predicate that describes the exact set of states s such that when program P is started in s, if it terminates it will terminate in a state satisfying condition B."



$\{10 < y\}$

```
y := y + 1;
z := x + y;
```

$\{x < z\}$

Exercise: Give “weakest” preconditions

1. $\{? \qquad\quad\}$ `x := x + 2` $\{x \geq 5\}$

Exercise: Give “weakest” preconditions

1. $\{\ x \geq 3\}$ `x := x + 2` $\{x \geq 5\}$

2.

$\{?\qquad\qquad\qquad\qquad\qquad\qquad\}$

`if (y < 0) then x := x+1 else x := y`

$\{x > 0\}$

Exercise: Give "weakest" preconditions

1. $\{\ x \geq 3\}$ `x := x + 2` $\{x \geq 5\}$

2. $\{\ (y < 0 \wedge x > -1) \vee (y > 0)\}$
   `if (y < 0) then x := x+1 else x := y`
   $\{x > 0\}$

3. $\{? \qquad \}$ while $(x \leq 5)$ do x := x+1 $\{x = 6\}$

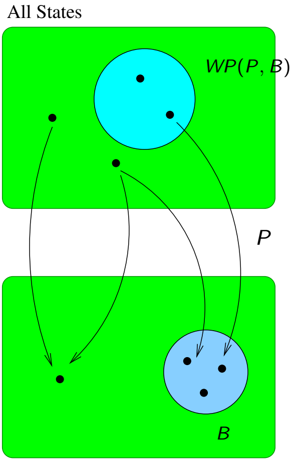
Exercise: Give “weakest” preconditions

1. $\{\ x \geq 3\}$ `x := x + 2` $\{x \geq 5\}$

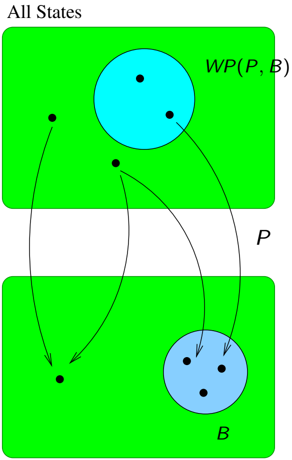
2. $\{\ (y < 0 \wedge x > -1) \vee (y > 0)\}$
   `if (y < 0) then x := x+1 else x := y`
   $\{x > 0\}$

3. $\{\ x \leq 6\}$ while $(x \leq 5)$ do x := x+1 $\{x = 6\}$

Exercise: How will you define $WP(P, B)$?



Exercise: How will you define $WP(P, B)$?



$$WP(P, B) = \{s \mid \forall t : (s, t) \in [P] \text{ we have } t \models B\}$$

Using weakest preconditions to partially automate inductive proofs

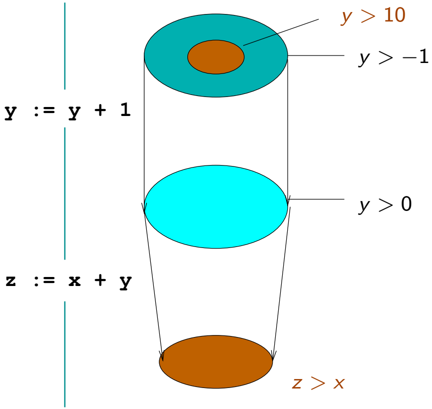
Weakest preconditions give us a way to:

- Check inductiveness of annotations

$$\{A_i\}\ S_i\ \{A_{i+1}\} \text{ iff } A_i \implies WP(S_i, A_{i+1})$$

- Reduce the amount of user-annotation needed
  - Programs without loops don't need any user-annotation
  - For programs with loops, user only needs to provide loop invariants

Checking $\{A\}\ P\ \{B\}$ using WP



Check that

$$(y > 10) \implies WP(P, z > x)$$

WP rules

- Hoare's rules for skip, assignment, and if-then-else are already WP rules.
- For Sequencing:

$$WP(S;T,\ B) = WP(S, WP(T, B)).$$

Weakest Precondition for `while` statements

- We can "approximate" $WP(\texttt{while}\ b\ \texttt{do}\ c)$.
- $WP_i(w, A) =$ the set of states from which the body c of the loop is either entered more than i times or we exit the loop in a state satisfying A.
- WP_i defined inductively as follows:

$$\begin{array}{lcl} WP_0 & = & b \vee A \\ WP_{i+1} & = & (\neg b \wedge A) \vee (b \wedge WP(c, WP_i)) \end{array}$$

- Then $WP(w, A)$ can be shown to be the "limit" or least upper bound of the chain $WP_0(w, A),\ WP_1(w, A), \ldots$ in a suitably defined lattice (here the join operation is "And" or intersection).

Illustration of WP_i through example

Consider the program w below:

`while` $(x \geq 10)$ `do`
    x := x - 1

- What is the weakest precondition of w with respect to the postcondition $(x \leq 0)$?
- Compute $WP_0(w, (x \leq 0))$, $WP_1(w, (x \leq 0))$, ...

Illustration of WP_i through example

Consider the program w below:

```
while (x ≥ 10) do
    x := x - 1
```

- What is the weakest precondition of w with respect to the postcondition $(x \leq 0)$?
- Compute $WP_0(w, (x \leq 0))$, $WP_1(w, (x \leq 0))$, ...

Automating checking of pre-post specifications for a program

To check:

$\{y > 10\}$

```
y := y + 1;
z := x + y;
```

$\{x < z\}$

Use the weakest precondition rules to generate the verification condition:

$$(y > 10) \implies (y > -1).$$

Check the verification condition by asking a theorem prover / SMT solver if the formula

$$(y > 10) \wedge \neg(y > -1).$$

is satisfiable.

What about `while` loops?

```
Pre: 0 <= n

int a := m;
int x := 0;
while (x < n) {
  a := a + 1;
  x := x + 1;
}

Post: a = m + n
```

Adequate loop invariant

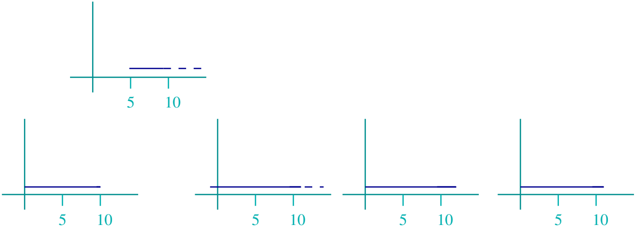
What is a “good” loop invariant for this program?

```
x := 0;
while (x < 10) {
 if (x >= 0)
  x := x + 1;
 else
  x := x - 1;
}
assert(x <= 11);
```

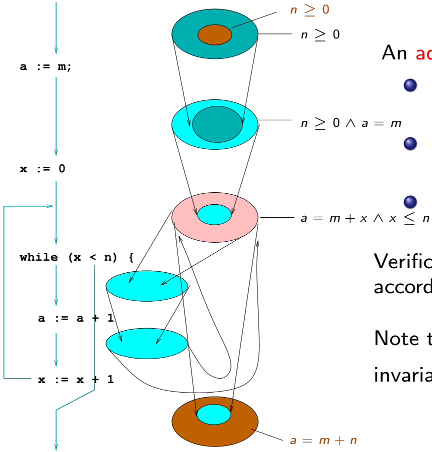
Adequate loop invariant

```
x := 0;
while (x < 10) {
 if (x >= 0)
  x := x + 1;
 else
  x := x - 1;
}
assert(x <= 11);
```

Canonical Invariant	*Not−inv*	*Inv,not−ind*	*Inv,ind,not−adeq*	*Inv,ind,adeq*
$0 \leq x \leq 10$	$5 \leq x$	$-1 \leq x$	$0 \leq x \leq 12$	$0 \leq x \leq 11$



Adequate loop invariant



An adequate loop invariant needs to satisfy:

- $\{n \geq 0\}$ `a := m; x := 0` $\{a = m + x \wedge x \leq n\}$.
- $\{a = m + x \wedge x \leq n \wedge x < n\}$ `a := a+1; x := x+1` $\{a = m + x \wedge x \leq n\}$.
- $\{a = m + x \wedge x \leq n \wedge x \geq n\}$ `skip` $\{a = m + n\}$.

Verification conditions are generated accordingly.

Note that $a = m + x$ is not an adequate loop invariant.

Generating Verification Conditions for a program

```
assume Pre

    S1

while (b) {        <-- invariant Inv

        S2

}

    S3

assert Post
```

The following VCs are generated:

- $Pre \wedge [S_1] \implies Inv'$
  Or: $Pre \implies WP(S_1, Inv)$
- $Inv \wedge b \wedge [S_2] \implies Inv'$
  Or: $(Inv \wedge b) \implies WP(S_2, Inv)$
- $Inv \wedge \neg b \wedge [S_3] \implies Post'$
  Or: $Inv \wedge \neg b \implies WP(S_3, Post)$

Relative completeness of Hoare logic

Theorem (Cook 1974)

Hoare logic is complete provided the assertion language L can express the WP for any program P and post-condition B.

Proof uses WP predicates and proceeds by induction on the structure of the program P.

- Suppose $\{A\}$ `skip` $\{B\}$ holds. Then it must be the case that $A \implies B$ is true. By Skip rule we know that $\{B\}$ `skip` $\{B\}$. Hence by Weakening rule, we get that $\{A\}$ `skip` $\{B\}$ holds.
- Suppose $\{A\}$ `x := e` $\{B\}$ holds. Then it must be the case that $A \implies B[e/x]$. By Assignment rule we know that $\{B[e/x]\}$ `x := e` $\{B\}$ is true. Hence by Weakening rule, we get that $\{A\}$ `x := e` $\{B\}$ holds.
- Similarly for sequencing `S;T`.
- Similarly for `if-then-else`.

Relative completeness of Hoare logic

- Suppose $\{A\}$ `while b do S` $\{B\}$ holds. Let $P = WP($`while b do S`$, B)$. Then it is not difficult to check that P is a loop invariant for the while statement. I.e $\{P \wedge b\}$ `S` $\{P\}$ is true. By induction hypothesis, this triple must be provable in Hoare logic. Hence we can conclude using the While rule, that $\{P\}$ `while b do S` $\{P \wedge \neg b\}$. But since P was a valid precondition, it follows that $(P \wedge \neg b) \implies B$. For the same reason, we should have $A \implies P$. By the weakening rule, we have a proof of $\{A\}$ `while b do S` $\{B\}$.

Conclusion

- Features of this Floyd-Hoare style of verification:
  - Tries to find a proof in the form of an inductive annotation.
  - A Floyd-style proof can be used to obtain a Hoare-style proof; and vice-versa.
  - Reduces verification (given key annotations) to checking satisfiability of a logical formula (VCs).
  - Is flexible about predicates, logic used (for example can add quantifiers to reason about arrays).
- Main challenge is the need for user annotation (adequate loop invariants).
- Can be increasingly automated (using learning techniques).