Interprocedural Analysis: Sharir-Pnueli's Call-strings Approach

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Outline

Handling programs with procedure calls

How would we extend an abstract interpretation to handle
programs with procedures?

```
main(){              f(){                 g(){
  x := 0;              x := x+1;            f();
  f();                 return;              return;
  g();               }                    }
  print x;
}
```

Handling programs with procedure calls

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Question: what is the collecting state before the print x statement in main?

Handling programs with procedure calls

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programs with procedures?

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main(){              f(){                 g(){
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  print x;
}
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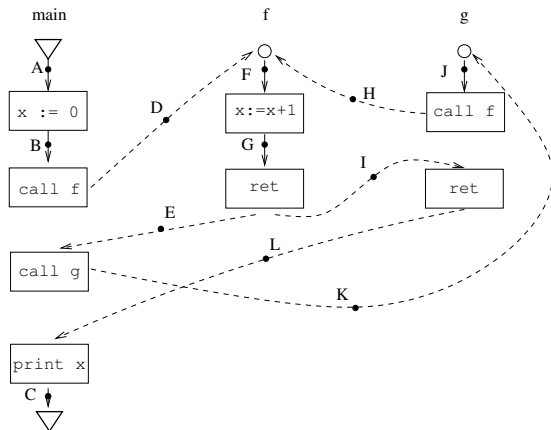
Question: what is the collecting state before the print x
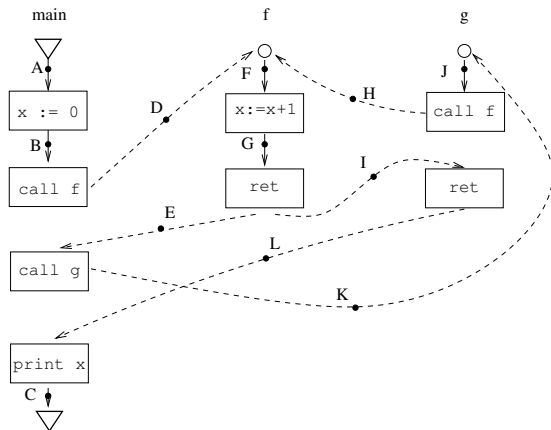statement in main? Answer: $x \mapsto 2$.

Handling programs with procedure calls

- Add extra edges
  - call edges: from call site (call p) to start of procedure (p)
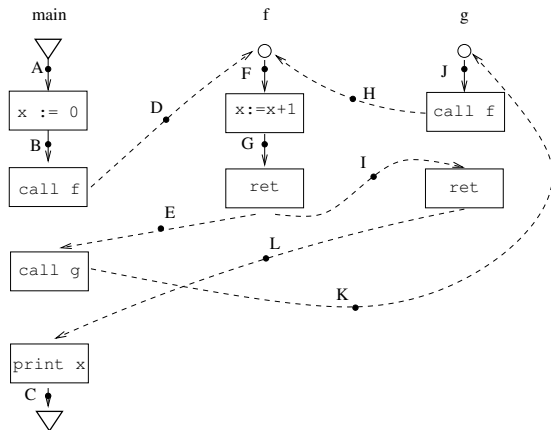 - ret edges: from return statement (in p) to point after call sites ("ret sites") (call p).

Handling programs with procedure calls

- Assume variables are uniquely named across program.
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Handling programs with procedure calls

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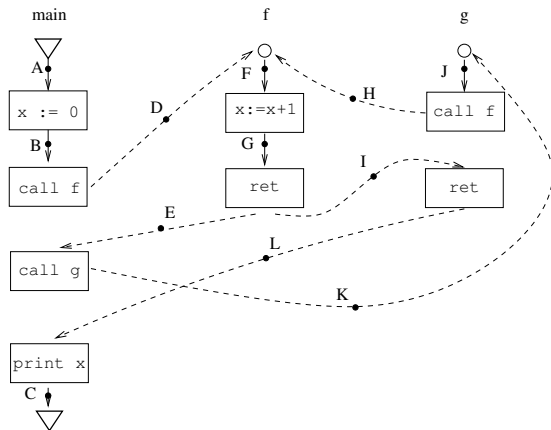
Handling programs with procedure calls

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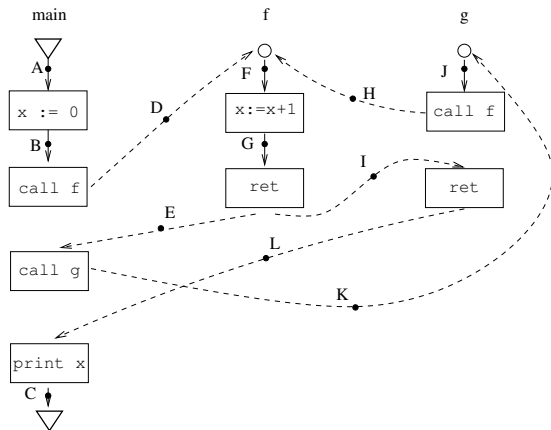- Now compute JOP in this extended control-flow graph.
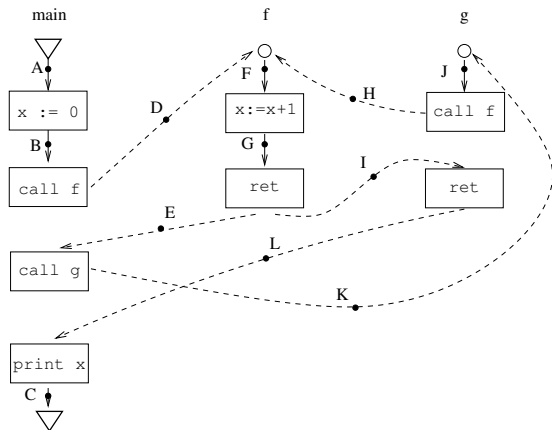
Problem with JOP in this graph

Ex. 1. Actual collecting
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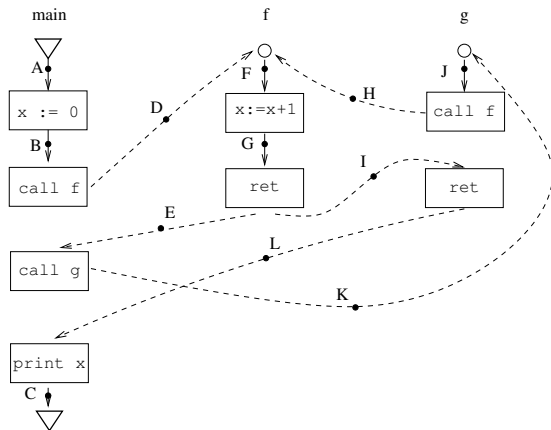
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Problem with JOP in this graph

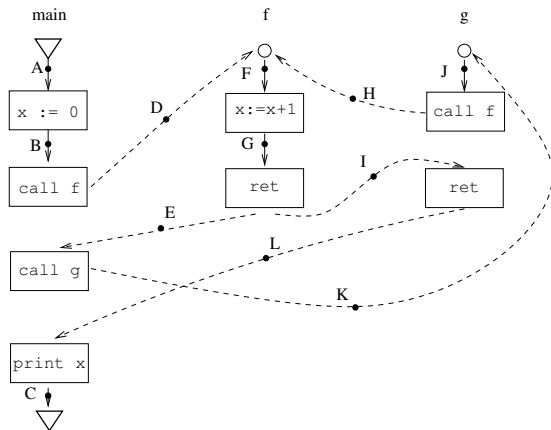Ex. 1. Actual collecting state at C? $\{x \mapsto 2\}$.
Ex. 2. JOP at C using collecting analysis?

Problem with JOP in this graph

Ex. 1. Actual collecting state at C? $\{x \mapsto 2\}$.
Ex. 2. JOP at C using collecting analysis? $\{x \mapsto 1,\ x \mapsto 2,\ x \mapsto 3, \ldots\}$.

Problem with JOP in this graph

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$\{x \mapsto 1, \ x \mapsto 2, \ x \mapsto 3, \ldots\}$.

- JOP is sound but
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- Reason: Some
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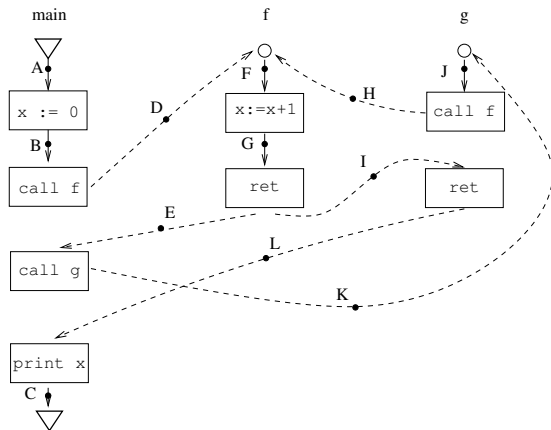 correspond to
  executions of the
  program: Eg.
  ABDFGILC.

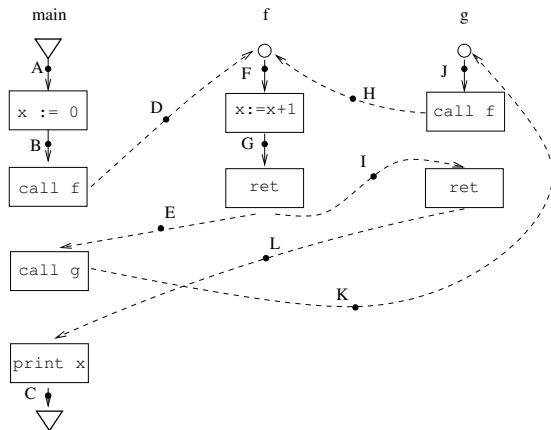## Problem with JOP in this graph

Ex. 1. Actual collecting state at C? $\{x \mapsto 2\}$.

Ex. 2. JOP at C using collecting analysis? $\{x \mapsto 1, \ x \mapsto 2, \ x \mapsto 3, \dots\}$.

- JOP is sound but very imprecise.

- Reason: Some paths don't correspond to executions of the program: Eg. ABDFGILC.

What we want is Join over "Interprocedurally-Valid" Paths (JVP).

Interprocedurally valid paths and their call-strings

- Informally a path ρ in the extended CFG G' is inter-procedurally valid if every return edge in ρ "corresponds" to the most recent "pending" call edge.
- For example, in the example program the ret edge E corresponds to the call edge D.
- The call-string of a valid path ρ is a subsequence of call edges which have not been "returned" as yet in ρ.
- For example, $cs(ABDFGEKJHF)$ is "KH".

Interprocedurally valid paths and their call-strings

- A path $\rho = ABDFGEKJHF$ in $IVP_{G'}$ for example program:
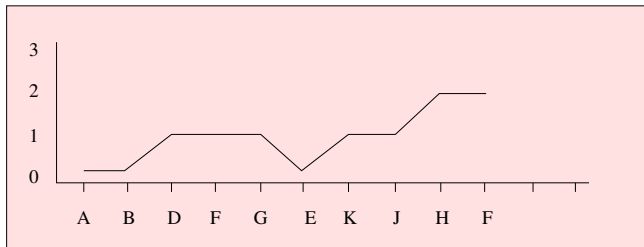
- Associated call-string $cs(\rho)$ is KH.
- For $\rho = ABDFGEK$ $cs(\rho) = K$.
- For $\rho = ABDFGE$ $cs(\rho) = \epsilon$.

Sharir and Pnueli's approaches to interprocedural analysis

Micha Sharir and Amir Pnueli: Two approaches to interprocedural data flow analysis, in *Program Flow Analysis: Theory and Applications* (Eds. Muchnick and Jones) (1981).

Interprocedurally valid paths and their call-strings

More formally: Let ρ be a path in G'. We define when ρ is interprocedurally valid (and we say $\rho \in IVP(G')$) and what is its call-string $cs(\rho)$, by induction on the length of ρ.

- If $\rho = \epsilon$ then $\rho \in IVP(G')$. In this case $cs(\rho) = \epsilon$.
- If $\rho = \rho' \cdot N$ then $\rho \in IVP(G')$ iff $\rho' \in IVP(G')$ with $cs(\rho') = \gamma$ say, and one of the following holds:
  1. N is neither a call nor a ret edge.
     In this case $cs(\rho) = \gamma$.
  2. N is a call edge.
     In this case $cs(\rho) = \gamma \cdot N$.
  3. N is ret edge, and γ is of the form $\gamma' \cdot C$, and N corresponds to the call edge C.
     In this case $cs(\rho) = \gamma'$.
- We denote the set of (potential) call-strings in G' by Γ. Thus $\Gamma = \mathcal{C}^*$, where $\mathcal{C}$ is the set of call edges in G'.

Join over interprocedurally-valid paths (JVP)

- Let P be a given program, with extended CFG G'.
- Let $path_{I,N}(G')$ be the set of paths from the initial point I to point N in G'.
- Let $\mathcal{A} = ((D, \leq), f_{MN}, d_0)$ be a given abstract interpretation.
- Then we define the join over all interprocedurally valid paths (JVP) at point N in G' to be:

$$\bigsqcup_{\rho \, \in \, path_{I,N}(G') \cap IVP(G')} f_\rho(d_0).$$

One approach to obtain JVP

- Find JOP over same graph, but modify the abs int.
- Modify transfer functions for call/ret edges to detect and invalidate invalid edges.
- Augment underlying data values with some information for this.
- Natural thing to try: "call-strings".
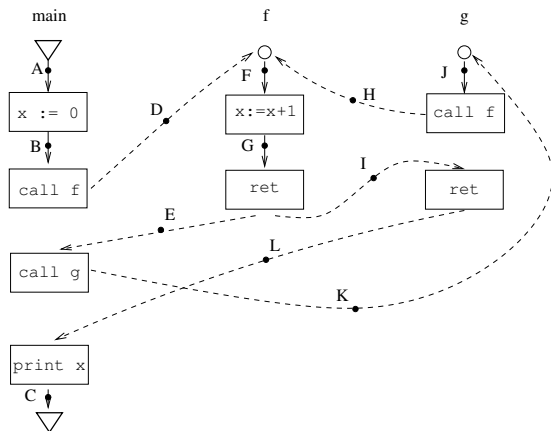
Overall plan

- Define an abs int $\mathcal{A}'$ which extends given abs int $\mathcal{A}$ with call-string data.
- Show that JOP of $\mathcal{A}'$ on G' coincides with JVP of $\mathcal{A}$ on G'.
- Use Kildall (or any other technique) to compute LFP of $\mathcal{A}'$ on G'. This value over-approximates JVP of $\mathcal{A}$ on G'.
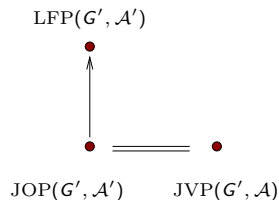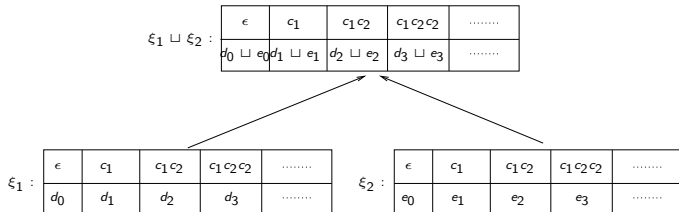
$\mathrm{LFP}(G', \mathcal{A}')$

$\mathrm{JOP}(G', \mathcal{A}')$ $\quad$ $\mathrm{JVP}(G', \mathcal{A})$

Call-string abs int $\mathcal{A}'$: Lattice $(D', \leq')$

- Elements of D' are maps $\xi : \Gamma \to D$

$\xi :$	ϵ	c_1	$c_1 c_2$	$c_1 c_2 c_2$	....
	d_0	d_1	d_2	d_3	....

- Ordering on D': $\leq'$ is the pointwise extension of $\leq$ in D.
- That is $\xi_1 \leq' \xi_2$ iff for each $\gamma \in \Gamma$, $\xi_1(\gamma) \leq \xi_2(\gamma)$.

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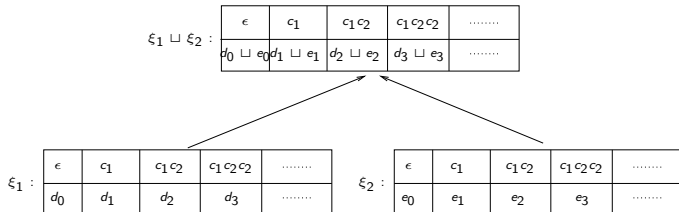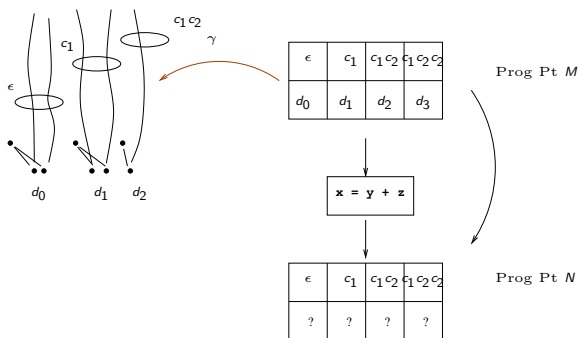
$\xi_1 \sqcup \xi_2 :$	ϵ	c_1	$c_1 c_2$	$c_1 c_2 c_2$	....
	$d_0 \sqcup e_0$	$d_1 \sqcup e_1$	$d_2 \sqcup e_2$	$d_3 \sqcup e_3$	....

$\xi_1 :$	ϵ	c_1	$c_1 c_2$	$c_1 c_2 c_2$	....
	d_0	d_1	d_2	d_3	....

$\xi_2 :$	ϵ	c_1	$c_1 c_2$	$c_1 c_2 c_2$	....
	e_0	e_1	e_2	e_3	....

- Check that $(D', \leq')$ is also a complete lattice.

Meaning of abstract values in $\mathcal{A}'$

- A call-string table ξ at program point N represents the fact that, for each call-string γ, there are some initial paths with call-string γ reaching N, and the join of the abstract states (obtained by propagating d_0) along these paths is $\xi(\gamma)$.
- The transfer functions of $\mathcal{A}'$ should keep this meaning in mind.

Call-string abs int $\mathcal{A}'$: Initial value ξ_0

- Initial value ξ_0 is given by

$$\xi_0(\gamma) = \left\{ \begin{array}{ll} d_0 & \text{if } \gamma = \epsilon \\ \bot & \text{otherwise.} \end{array} \right.$$

$\xi_0:$

ϵ	c_1	$c_1 c_2$	$c_1 c_2 c_2$	....
d_0	$\bot$	$\bot$	$\bot$	....

Call-string abs int $\mathcal{A}'$: transfer functions

- Transfer functions for non-call/ret edge N:

$$f'_{MN}(\xi) = f_{MN} \circ \xi.$$

- Transfer functions for call edge N:

$$f'_{MN}(\xi) = \lambda\gamma. \begin{cases} \xi(\gamma') & \text{if } \gamma = \gamma' \cdot N \\ \bot & \text{otherwise} \end{cases}$$

- Transfer functions for ret edge N whose corresponding call edge is C:

$$f'_{MN}(\xi) = \lambda\gamma.\xi(\gamma \cdot C)$$

- Transfer functions f'_{MN} is monotonic (distributive) if each f_{MN} is monotonic (distributive).

Transfer functions f'_{MN} for example program

- Non-call/ret edge B:

  $$\xi_B = f_{AB} \circ \xi_A.$$

- Call edge D:

  $$\xi_D(\gamma) = \begin{cases} \xi_B(\gamma') & \text{if } \gamma = \gamma' \cdot D \\ \bot & \text{otherwise} \end{cases}$$

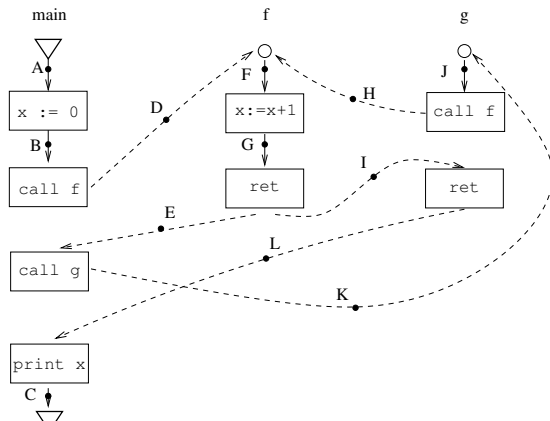- Return edge E:

  $$\xi_E(\gamma) = \xi_G(\gamma \cdot D).$$
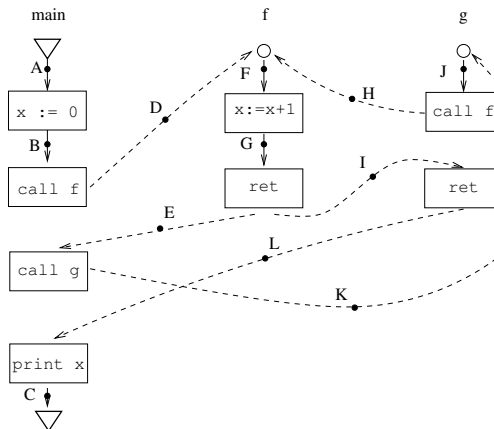
Exercise 1

Let $\mathcal{A}$ be the standard collecting state analysis. For brevity, represent a set of concrete states as $\{0, 1\}$ (meaning the 2 concrete states $x \mapsto 0$ and $x \mapsto 1$). Assume an initial value $d_0 = \{0\}$.

Show the call-string tagged abstract states (in the lattice $\mathcal{A}'$) along the paths

1. ABDFGEKJHFGIL (interprocedurally valid)
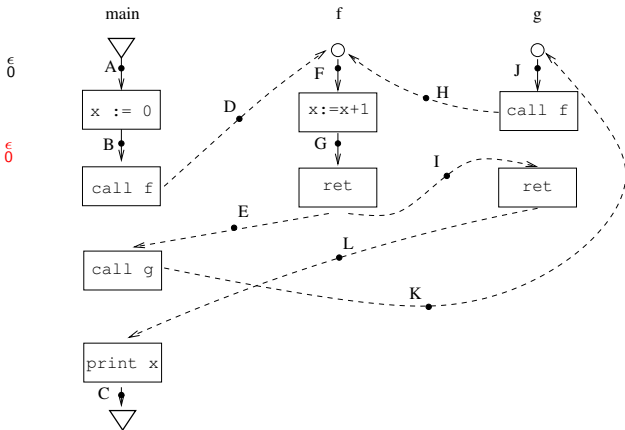2. ABDFGIL (interprocedurally invalid).

Exercise 2

Use Kildall's algo to compute the LFP of the $\mathcal{A}'$ analysis for the example program. Start with initial value $d_0 = \{0\}$.
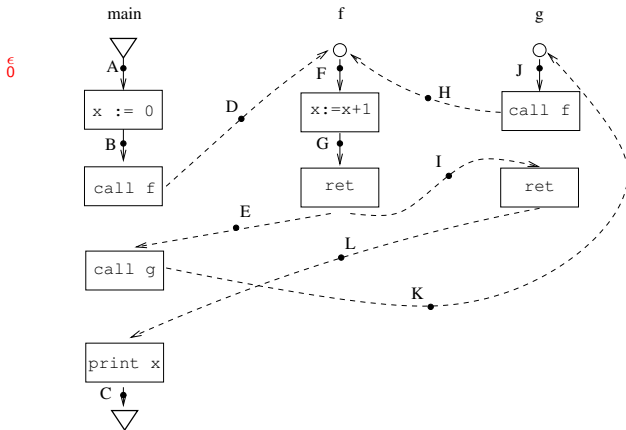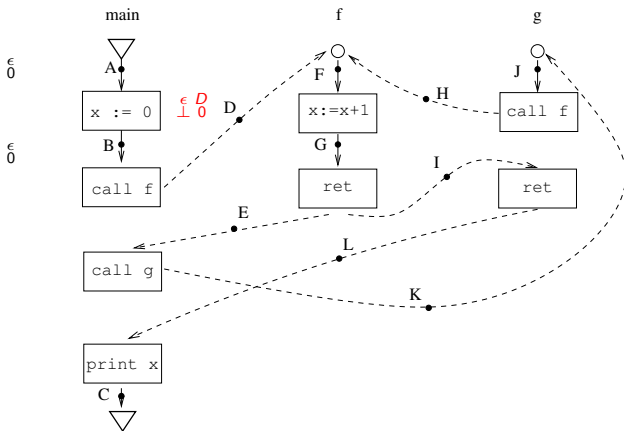
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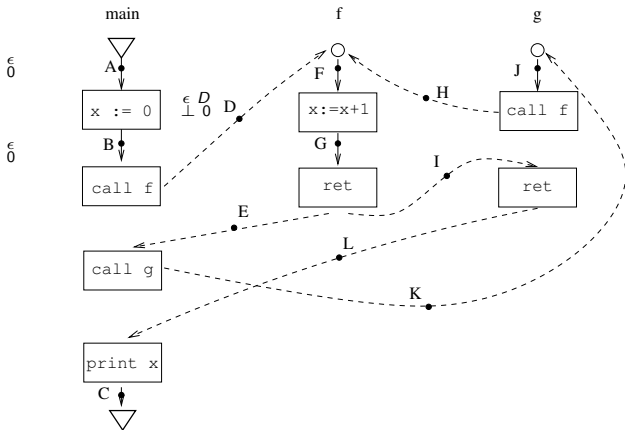
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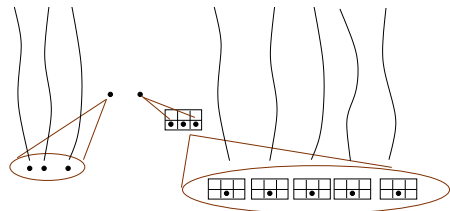
Correctness claim

Assumption on $\mathcal{A}$: Each transfer function satisfies $f_{MN}(\bot) = \bot$.

Claim

Let N be a point in G'. Then

$$JVP_{\mathcal{A}}(N) = \bigsqcup_{\gamma \in \Gamma} JOP_{\mathcal{A}'}(N)(\gamma).$$

Proof: Use following lemmas to prove that LHS dominates RHS and vice-versa.

IVP Paths reaching N   Paths reaching N

Correctness claim: Lemma 1

Lemma 1

Let ρ be a path in $IVP_{G'}$. Then

$$f'_\rho(\xi_0) = \lambda\gamma. \left\{ \begin{array}{ll} f_\rho(d_0) & if\, \gamma = cs(\rho) \\ \bot & \text{otherwise.} \end{array} \right.$$

ϵ	c_1	$cs(\rho)$	$c_1 c_2 c_2$	....
$\bot$	$\bot$	d	$\bot$	....

Proof: by induction on the length of ρ.

Correctness claim: Lemma 2

Lemma 2

Let ρ be a path not in $IVP_{G'}$. Then

$$f'_\rho(\xi_0) = \lambda\gamma.\bot.$$

ϵ	c_1	c_2	$c_1 c_2 c_2$	....
$\bot$	$\bot$	$\bot$	$\bot$	....

Proof:

- ρ must have an invalid prefix.
- Consider smallest such prefix $\alpha \cdot N$. Then it must be that α is valid and N is a return edge not corresponding to $cs(\alpha)$.
- Using previous lemma it follows that $f'_{\alpha \cdot N}(\xi_0) = \lambda\gamma.\bot.$
- But then all extensions of α along ρ must also have transfer function $\lambda\gamma.\bot.$

Computing JOP for abs int $\mathcal{A}'$

- Problem is that D' is infinite in general (even if D were finite). So we cannot use Kildall's algo to compute an over-approximation of JOP.
- We give two methods to bound the number of call-strings
  - Use "approximate" call-strings.
  - Give a bound on largest call-string needed.

Approximate (suffix) call-string method

Idea:

- Consider only call-strings of up to length $\leq l$, that may additionally be prefixed by a "$*$".
- A "$*$" prefix means that we have left out some initial calls.
- For $l = 2$, call strings can be of the form "$c_1 c_2$" or "$* c_1 c_2$" etc. So each table ξ is now a finite table.
- Transfer functions for non-call/ret edges remain same.
- Transfer functions for call edge C: Shift γ entry to $\gamma \cdot C$ if $|\gamma \cdot C| \leq l$; else shift it to $* \cdot \gamma' \cdot C$ where γ is of the form $A \cdot \gamma'$, for some call A.
- Transfer functions for ret edge N:
  - If $\gamma = \gamma' \cdot C$ and N corresponds to call edge C, then shift $\gamma' \cdot C$ entry to γ' entry.
  - If $\gamma = *$ then copy its entry to $*$ entry at the return site.
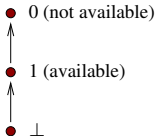
Available expressions analysis

- 0 (not available)
- 1 (available)
- ⊥
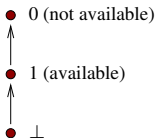
Lattice for Av-Exp analysis.

- Is a*b available at program point *N*?

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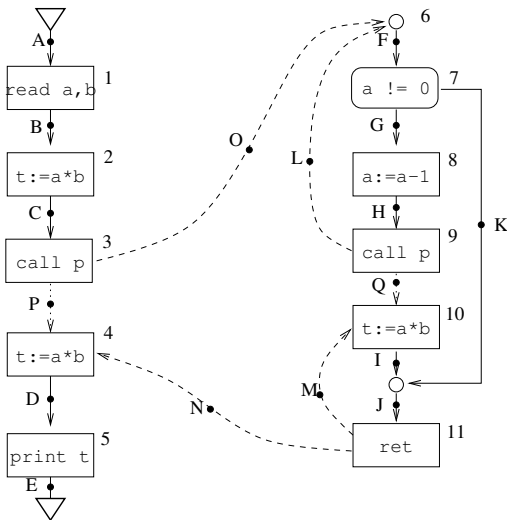
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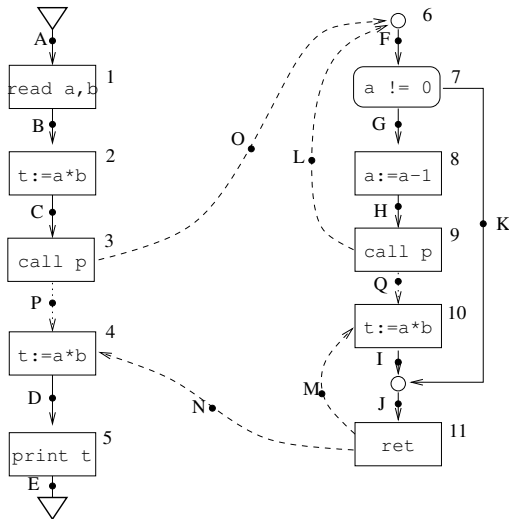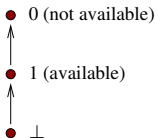- 1 (available)
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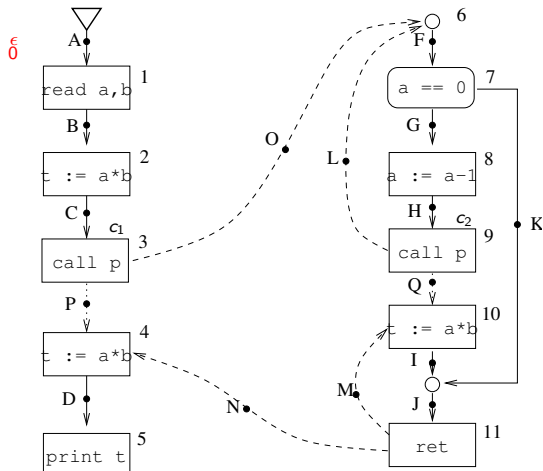## Available expressions analysis

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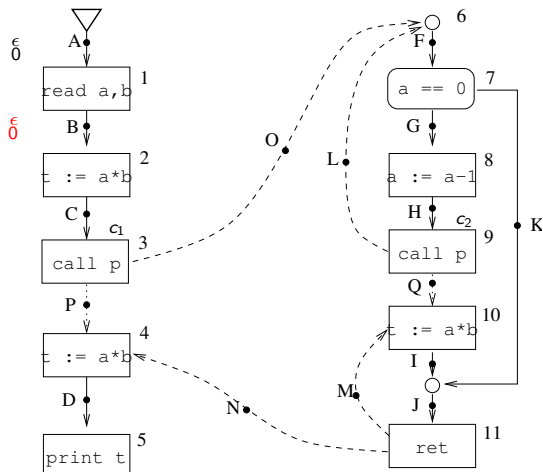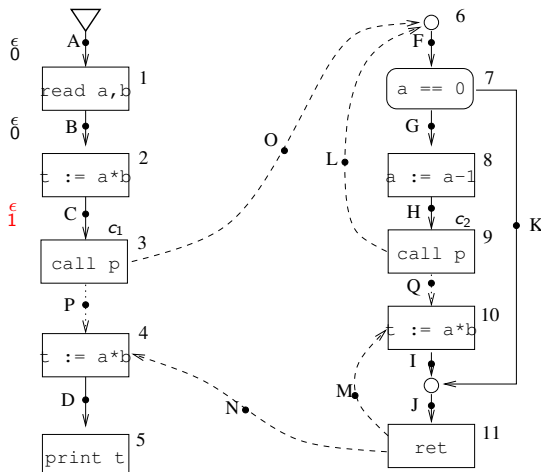- Yes if we consider interprocedurally valid paths only.

Exercise: approximate call-strings

Assume approximate call-string length of 2. Use Kildall's algo to compute the ξ table values for the example program. Start with initial value $d_0 = 0$.

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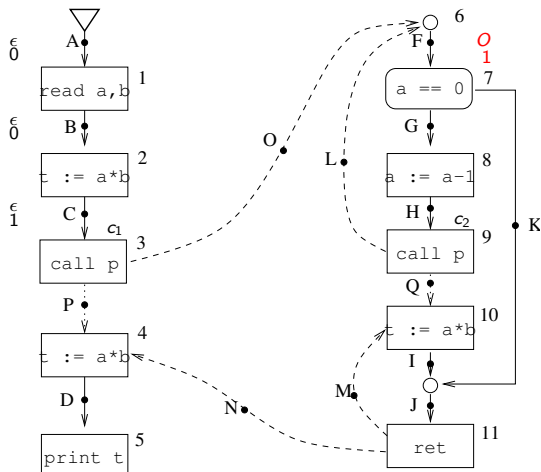
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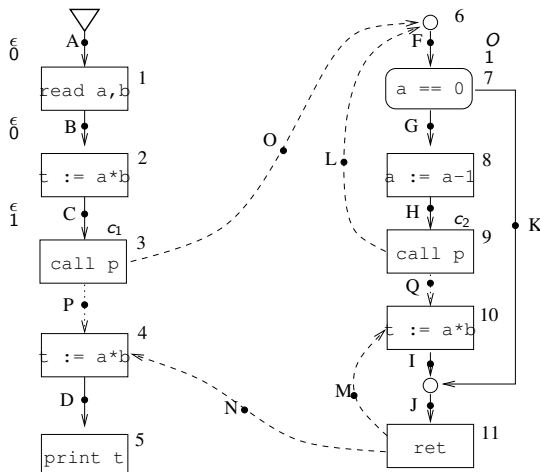
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Bounded call-string method for finite underlying lattice D

- Possible to bound length of call-strings Γ we need to consider.
- For a number l, we denote the set of call-strings (for the given program P) of length at most l, by Γ_l.
- Define a new analysis $\mathcal{A}''$ (M-bounded call-string analysis) in which call-string tables have entries only for Γ_M for a certain constant M, and transfer functions ignore entries for call-strings of length more than M.
- We will show that $\mathrm{JOP}(G', \mathcal{A}'') = \mathrm{JOP}(G', \mathcal{A}')$.

$$\mathrm{JOP}(G', \mathcal{A}'') \quad \mathrm{JOP}(G', \mathcal{A}') \quad \mathrm{JVP}(G', \mathcal{A})$$
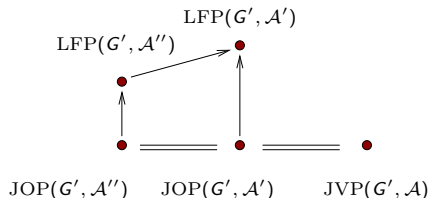
LFP of $\mathcal{A}''$ is more precise than LFP of $\mathcal{A}'$

- Consider any fixpoint V' (a vector of tables) of $\mathcal{A}'$.
- Truncate each entry of V' to (call-strings of) length M, to get V''.
- Clearly V' dominates V''.
- Further, observe that V'' is a post-fixpoint of the transfer functions for $\mathcal{A}''$.
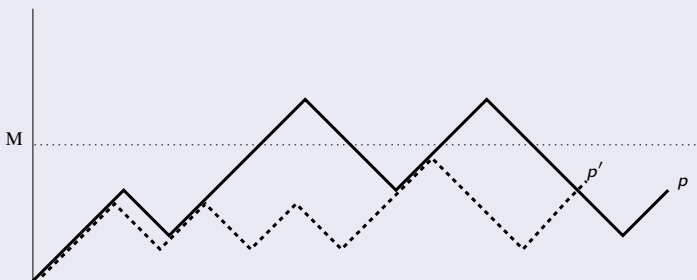- By Knaster-Tarski characterisation of LFP, we know that V'' dominates $\mathrm{LFP}(\mathcal{A}'')$.

$$\mathrm{JOP}(G', \mathcal{A}'') \quad \mathrm{JOP}(G', \mathcal{A}') \quad \mathrm{JVP}(G', \mathcal{A})$$
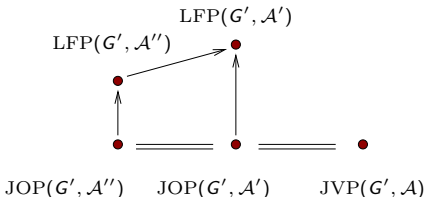
Sufficiency (or safety) of bound

Let k be the number of call sites in P.

Claim

For any path p in $IVP(r_1, N)$ with a prefix q such that $|cs(q)| > k|D|^2 = M$ there is a path p' in $IVP(r_1, N)$ with $|cs(q')| \leq M$ for each prefix q' of p', and $f_p(d_0) = f_{p'}(d_0)$.

Paths with bounded call-strings

Proving claim

Claim

For any path p in $IVP(r_1, N)$ such that for some prefix q of p, $|cs(q)| > M = k|D|^2$, there is a path p' in $IVP_{\Gamma_M}(r_1, N)$ with $f_{p'}(d_0) = f_p(d_0)$.

- Sufficient to prove:

Subclaim

For any path p in $IVP(r_1, N)$ with a prefix q such that $|cs(q)| > M$, we can produce a smaller path p' in $IVP(r_1, N)$ with $f_{p'}(d_0) = f_p(d_0)$.
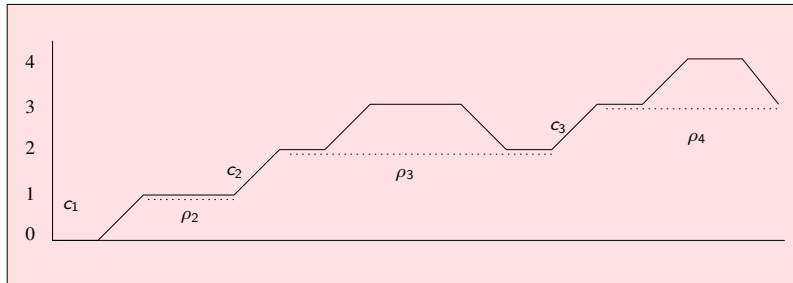

- ...since if $|p| \leq M$ then $p \in IVP_{\Gamma_M}$.

Proving subclaim: Path decomposition

A path ρ in $IVP(r_1, n)$ can be decomposed as

$$\rho_1 \| (c_1, r_{p_2}) \| \rho_2 \| (c_2, r_{p_3}) \| \sigma_3 \| \cdots \| (c_{j-1}, r_{p_j}) \| \rho_j.$$

where each ρ_i $(i < j)$ is a valid and complete path from r_{p_i} to c_i, and ρ_j is a valid and complete path from r_{p_j} to n. Thus $c_1, \ldots, c_{j-1}$ are the unfinished calls at the end of ρ.
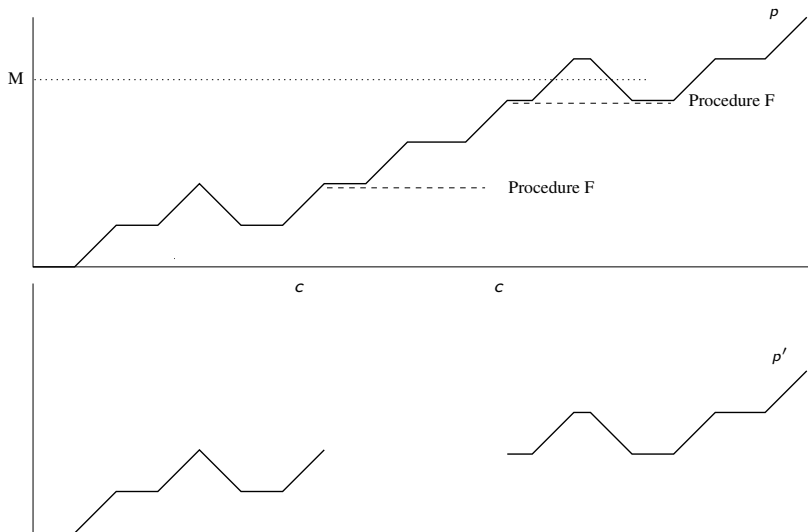
Proving subclaim

- Let p_0 be the first prefix of p where $|cs(p_0)| > M$.
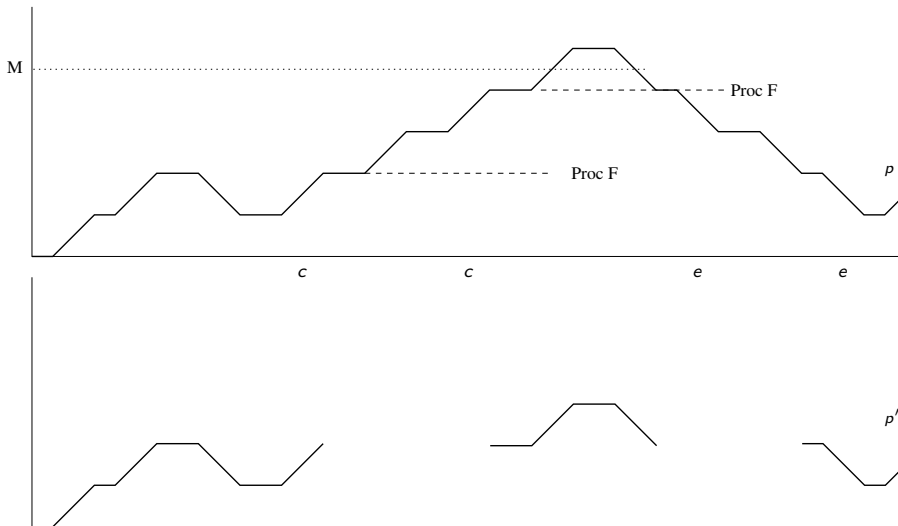- Let decomposition of p_0 be

$$\rho_1 \| (c_1, r_{p_2}) \| \rho_2 \| (c_2, r_{p_3}) \| \sigma_3 \| \cdots \| (c_{j-1}, r_{p_j}) \| \rho_j.$$

- Tag each unfinished-call c in p_0 by $(c, f_{q \cdot c}(d_0), f_{q \cdot cq'e}(d_0))$ where e is corresponding return of c in p.
- If no return for c in p tag with $(c, f_{q \cdot c}(d_0), \bot)$.
- Number of distinct such tags is $k \cdot |D|^2$.
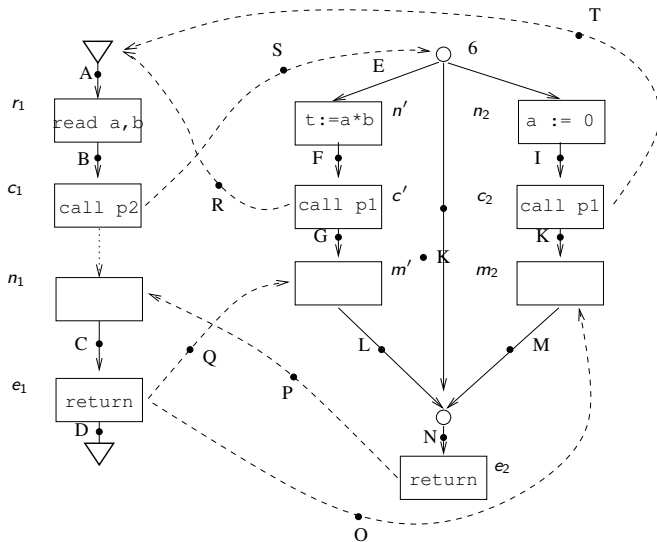- So there are two calls qc and $qcq'c$ with same tag values.

Proving subclaim – tag values are $\perp$

Proving subclaim – tag values are not $\perp$

Example

Transfer functions f'_{MN} for Example 2

- Non-call/ret edge C:

$$\xi_C = f_{BC} \circ \xi_B.$$

- Call edge O:

$$\xi_O(\gamma) = \begin{cases} \xi_C(\gamma') & \text{if } \gamma = \gamma' \cdot O_{\texttt{call p}} \\ \bot & \text{otherwise} \end{cases}$$

- Return edge N:

$$\xi_N(\gamma) = \xi_J(\gamma \cdot O).$$